Compilersupported ILP

STOP DOING Optimization

- Code was never meant to be optimized
- YEARS OF OPTIMIZING yet NO REAL-WORLD USE FOUND for BETTER PERFORMANCE
- Wanted to get better performance anyways? We had a tool for that, it was called "Upgrading hardware"
- "Yes please give me a low memory footprint. Please give me 5% CPU utilization" - Statements dreamed up by the utterly Deranged

Look at what Low-level programmers have been demanding your Respect for all this time, with all the RAM & CPU cores we built for them

(This is REAL optimizations, done by REAL programeers):





y = number; i = *(long *) &y; i = 0x5f3759df - (i >> 1); y = *(float *) &i; y = y *(threehalfs - (x2 * y * y));

???????

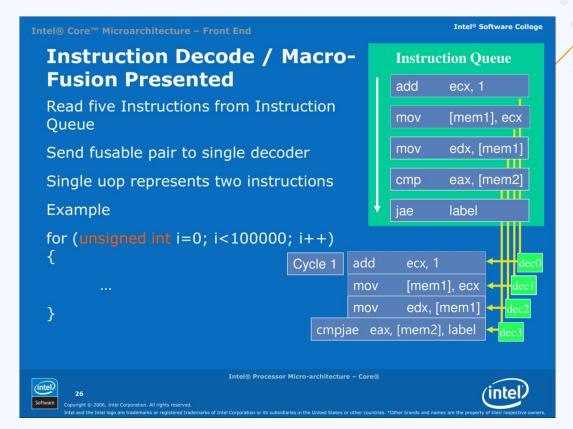
61 145 B B V V VIII A

"I spent the entire week reducing the system latency by 2ms"

They have played us for absolute fools

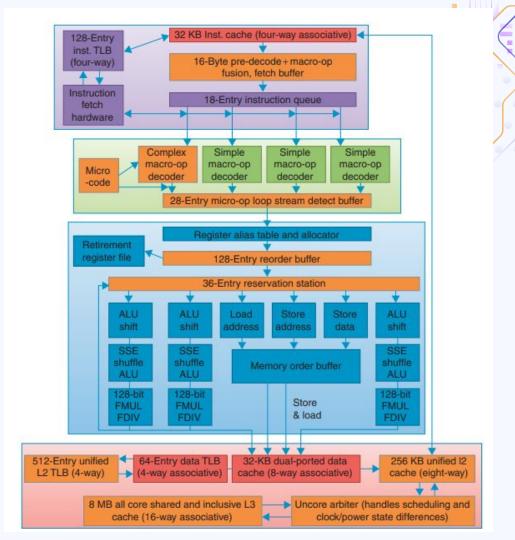
Macro-op fusion

source



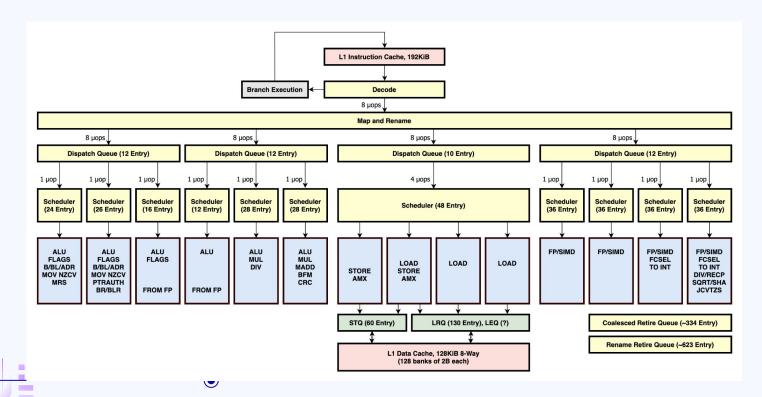
Intel Core i7 (H&P fig. 3.41)

Pre-decode??
Complex macro-op
decoder??
Loop stream detect??



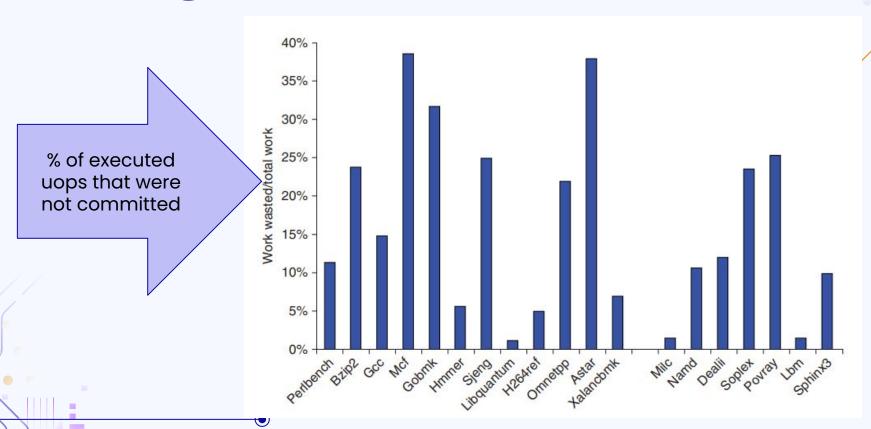
Firestorm (Apple M1)

Source (NOTE: reverse-engineered: might not be fully accurate)



What performance metrics (beyond CPI) might become important in a speculative CPU?

H&P fig. 3.42



What effects would speculative execution have on the memory system?

(Hint: think protected access and/or caches)

Remember this question?



What else could we parallelize here?

```
for (int i = 0; i < 100; i++) {
    A[i] = A[i] + B[i];
}</pre>
```

A simpler example

```
addi t0, x0, 0 // t0/i = 0
addi t1, x0, 100 // t1 = 100
loop: bge t0, t1, end
slli t2, t0, 2 // t2 = t0/i * 4
add t2, a0, t2 // t2 = A + t2
lw t3, 0(t2) // t3 = A[i]
add t3, t3, a1 // t3 = A[i] + c
sw t3, 0(t2) // A[i] = A[i] +c
addi t0, t0, 1 // t0/i++
j loop
end: nop
```

```
for (int i = 0; i < 100; i++) {
    A[i] = A[i] + c;
}</pre>
```

Reduce # of computations in loop

```
addi t0, x0, 0 // t0/i = 0
addi t1, \times 0, 100 // t1 = 100
loop: bge t0, t1, end
slli t2, t0, 2 // t2 = t0/i * 4
add t2, a0, t2 // t2 = A + t2
lw t3, O(t2) // t3 = A[i]
add t3, t3, a1 // t3 = A[i] + c
sw t3, 0(t2) // A[i] = A[i] +c
addi t0, t0, 1 // t0/i++
i loop
end: nop
```

```
addi t0, \times0, 0 // t0/i = 0
addi t1, x0, 100 // t1 = 100
addi t2, a0, 0 // t2 = A
loop: bge t0, t1, end
lw t3, 0(t2) // t3 = A[i]
add t3, t3, a1 // t3 = A[i] + c
sw t3, 0(t2) // A[i] = A[i] +c
addi t2, t2, 4 // advance pointer
addi t0, t0, 1 // t0/i++
j loop
end: nop
```

Get rid of i

```
addi t0, x0, 0 // t0/i = 0
addi t1, \times 0, 100 // t1 = 100
addi t2, a0, 0 // t2 = A
loop: bge t0, t1, end
lw t3, O(t2) // t3 = A[i]
add t3, t3, a1 // t3 = A[i] + c
sw t3, 0(t2) // A[i] = A[i] +c
addi t2, t2, 4 // advance pointer
addi t0, t0, 1 // t0/i++
j loop
end: nop
```

```
addi t2, a0, 0 // t2 = A

addi t1, a0, 400 // stop before A[100]

loop: bge t2, t1, end

lw t3, 0(t2) // t3 = A[i]

add t3, t3, a1 // t3 = A[i] + c

sw t3, 0(t2) // A[i] = A[i] +c

addi t2, t2, 4 // advance pointer

j loop
end: nop
```

What else could a compiler do?

Loop unrolling

3 3

What if we don't know the # of iterations through the loop at compile time?

int (i = 0; i < input_max; i++)

Dynamic loop size

```
for (int i = 0; i < input_max % 4; i++) {
      // do loop body
}
for (int j = input_max % 4; j < input_max; j += 4) {
      // unrolled loop body (4x)
}</pre>
```

Besides compiler complexity, what is a downside to loop unrolling?
(Could aggressive unrolling reduce performance?)

Could we unroll this loop?

```
for (int i = 0; i < 100; i++) {
    A[i + 1] = A[i] + C[i];
    B[i + 1] = B[i] + A[i + 1];
}</pre>
```

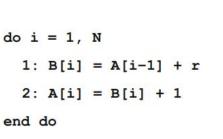
Could we unroll this loop?

```
for (int i = 0; i < 100; i++) {
    A[i] = A[i] + B[i];
    B[i + 1] = C[i] + D[i]
}</pre>
```

Detecting loop dependences

Figure 6.5: Intra-iteration and loop carried dependences

From *Instruction Level Parallelism* by Aiken, Banerjee, Kejariwal, Nicolau



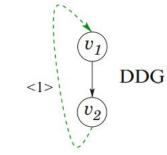


Figure 6.6 : Loop recurrences

Other static ILP approaches

Software pipelining

Pipeline dependent instructions within a loop

Global scheduling

Move instructions across basic blocks

Trace scheduling

Find a trace (common path through program) of multiple basic blocks

Rearrange and parallelize instructions within trace

Need "compensation code" in case branching into/out of trace

VLIW (Very Long Instruction Word)

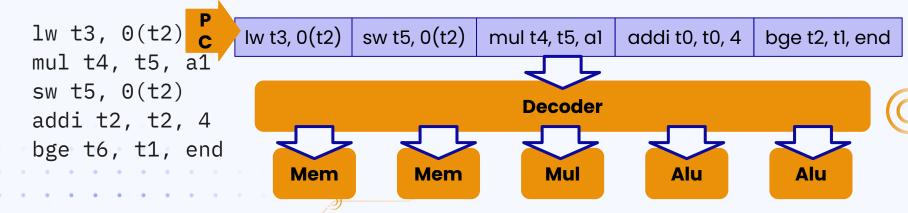
Compiler packs instructions into one long instruction word

Early VLIW: no dependences between instructions, units operate in lockstep

Pairs with loop unrolling, trace scheduling

Pros:

Cons:



What tradeoffs do you see between dynamic and static ILP? Which do you like more? Can you imagine ways to combine them?