

CSCI 1515 Applied Cryptography

This Lecture:

- Message Integrity
- Chosen-Message Attack (CMA) Security
- RSA Signature Scheme
- Authenticated Encryption

Message Integrity

Alice



"Let's meet @ 9am" →

Bob



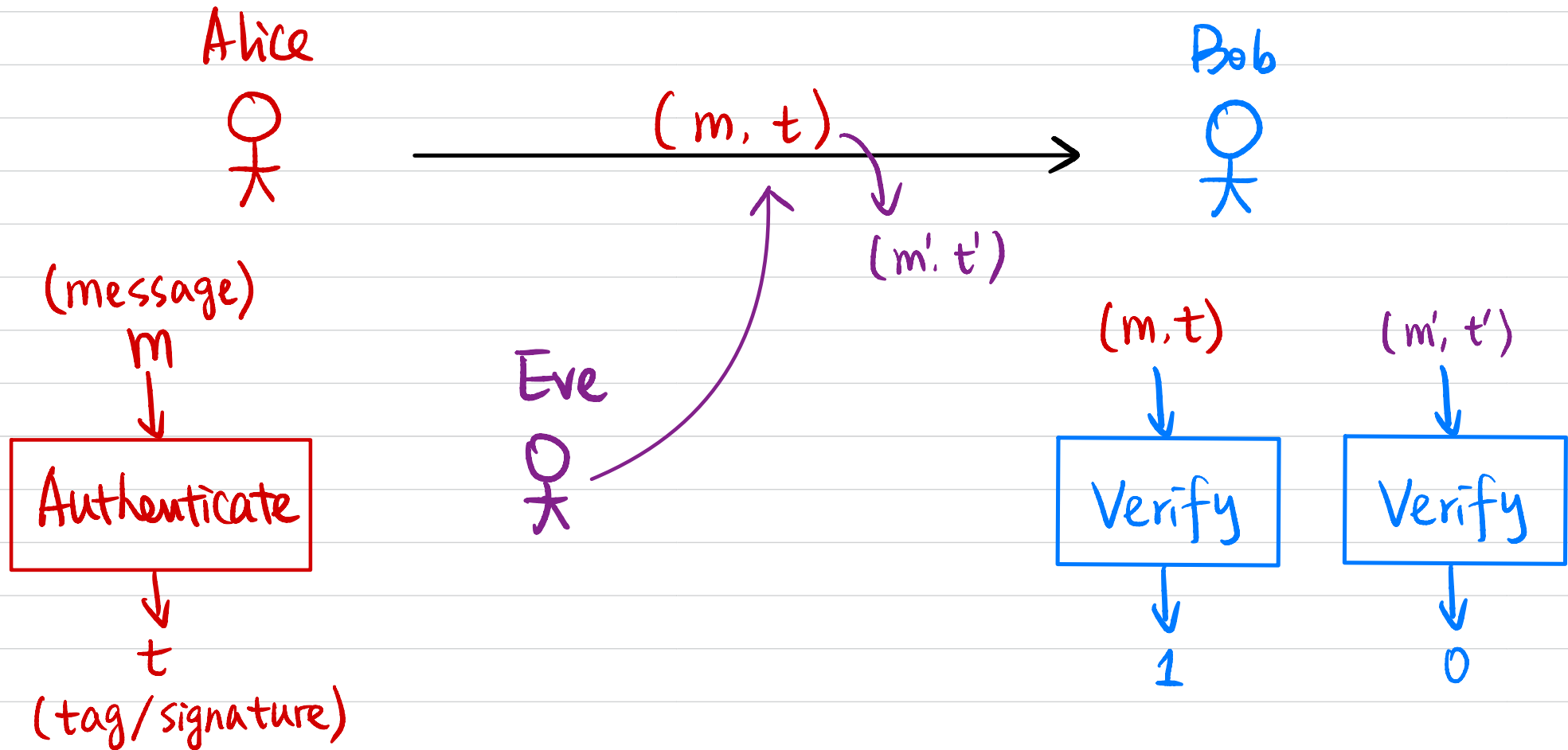
tamper with

Eve



Is it from Alice?

Message Integrity



Message Authentication Code (MAC)

Alice



(m, t)

Bob



(message)

m

k

Authenticate



t

(tag)



Eve



(m', t')

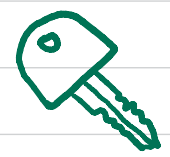
(m, t)

k

Verify



0/1



Digital Signature

Alice



(message)

m

sk

Authenticate



σ

(signature)



(secret)

(m, σ)

(m', σ')

Eve



Bob



(m, σ)

pk

Verify



0/1



(public)

Syntax

Message Authentication Code (MAC) Scheme $\Pi = (\text{Gen}, \text{Mac}, \text{Vrfy})$

$$k \leftarrow \text{Gen}(1^\lambda)$$

$$t \leftarrow \text{Mac}_k(m)$$

$$0/1 := \text{Vrfy}_k(m, t)$$

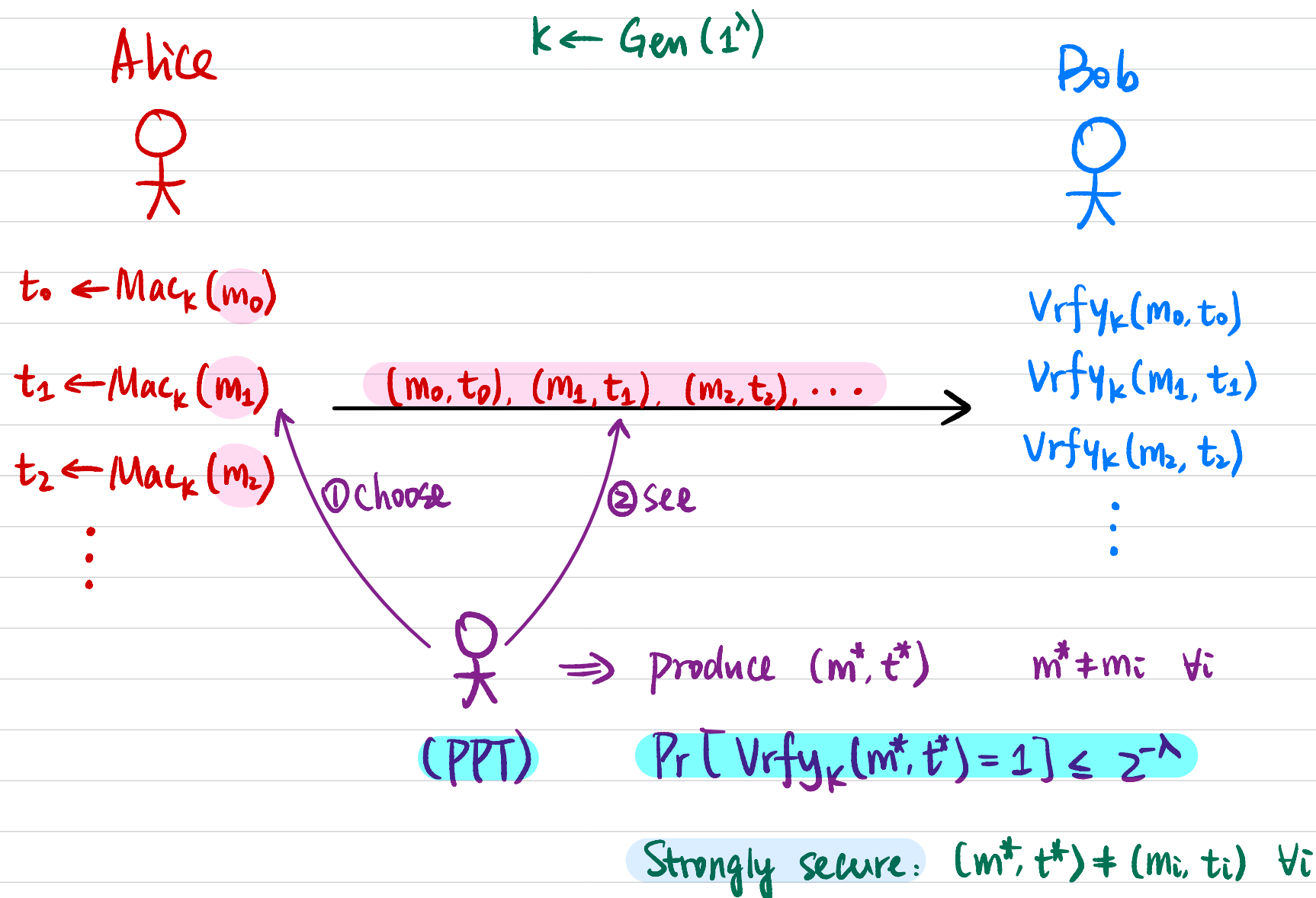
Digital Signature Scheme $\Pi = (\text{Gen}, \text{Sign}, \text{Vrfy})$

$$(pk, sk) \leftarrow \text{Gen}(1^\lambda)$$

$$\sigma \leftarrow \text{Sign}_{sk}(m)$$

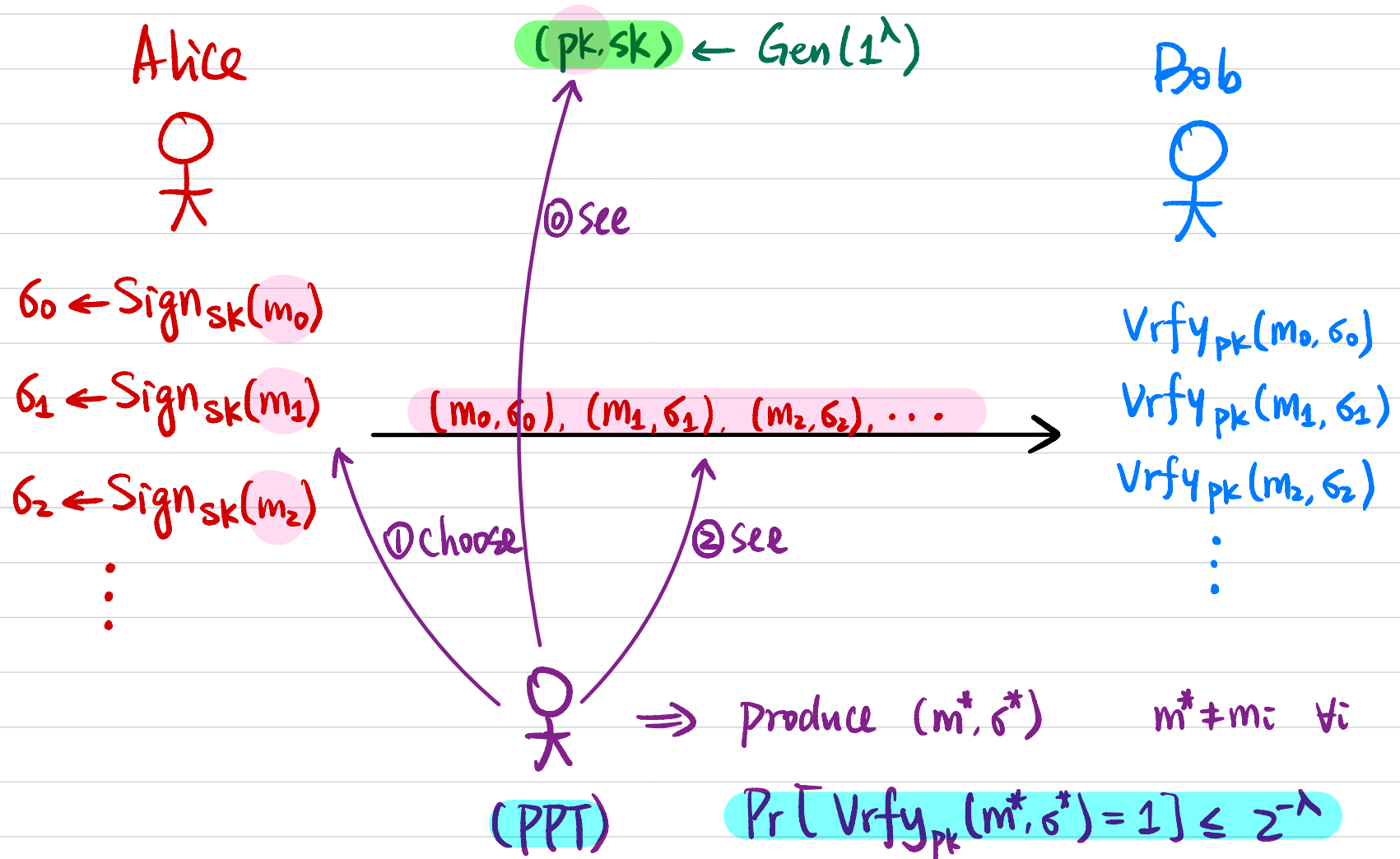
$$0/1 := \text{Vrfy}_{pk}(m, \sigma)$$

Chosen-Message Attack (CMA)



Does Mac have to be randomized?

Chosen-Message Attack (CMA)



Strongly secure: $(m_i^*, \sigma^*) \neq (m_i, \sigma_i) \forall i$

Does Sign have to be randomized?

RSA Signature

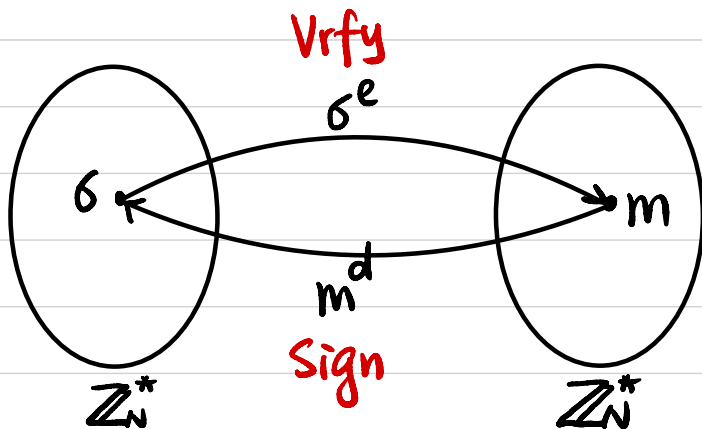
Generate two n -bit primes p, q ($p \neq q$)

Compute $N = p \cdot q$, $\phi(N) = (p-1)(q-1)$

Choose e s.t. $\gcd(e, \phi(N)) = 1$

Compute $d = e^{-1} \bmod \phi(N)$.

Given (N, e) & a random $y \leftarrow \mathbb{Z}_N^*$, it's computationally hard to find x s.t.
 $x^e \equiv y \pmod{N}$



$$sk = d \quad pk = (N, e)$$

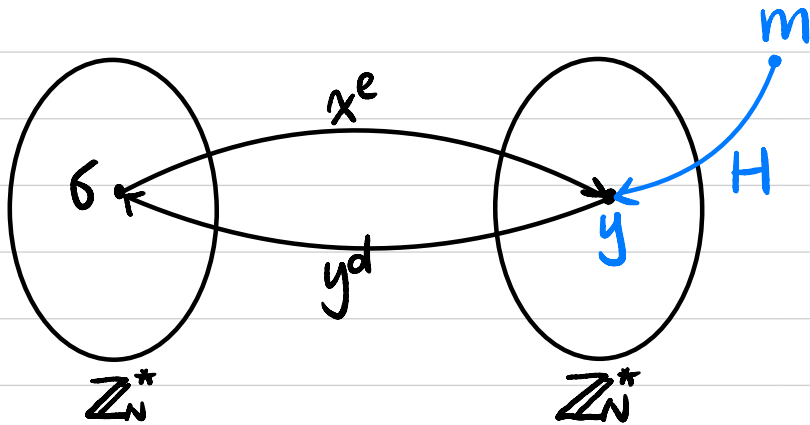
$$\text{Sign}_{sk}(m) = m^d \bmod N$$

$$\text{Vrfy}_{pk}(m, \sigma): \sigma^e \stackrel{?}{\equiv} m \pmod{N}$$

CMA Secure ?

RSA Signature

Given (N, e) & a random $y \leftarrow \mathbb{Z}_N^*$, it's computationally hard to find x s.t.
 $x^e \equiv y \pmod{N}$



$$sk = d \quad pk = (N, e)$$

$$\text{Sign}_{sk}(m) = H(m)^d \pmod{N}$$

$$\text{Vrfy}_{pk}(m, \sigma): \sigma^e \stackrel{?}{\equiv} H(m) \pmod{N}$$

$$H: \{0, 1\}^* \rightarrow \mathbb{Z}_N^*$$

↑
public, deterministic function
that gives a (pseudo)random output

CMA Secure ?

Constructions for Digital Signature

RSA Signature : RSA Assumption

Schnorr / DSA Signature : Discrete Logarithm Assumption

Lattice-Based Signature Schemes (Post-Quantum Security)

Constructions for MAC

From block cipher: CBC-MAC

From hash function: HMAC

Computational Assumption: "The construction is secure"

Summary

Symmetric-Key

Public-Key

Message
Secrecy

Primitive: SKE

Construction: block cipher

Primitive: PKE

Constructions: RSA / ElGamal

Message
Integrity

Primitive: MAC

Constructions: CBC-MAC / HMAC

Primitive: Signature

Constructions: RSA / DSA

Key Exchange

Construction: Diffie-Hellman

Important
Tool

Primitive: Hash function

Construction: SHA

What happens in practice

Alice



k

Diffie-Hellman Key Exchange

Bob



k

Symmetric-Key Encryption



Message Secrecy & Integrity?

Authenticated Encryption (AE)

Def An authenticated encryption scheme is a symmetric-key encryption that is CCA-secure & unforgeable.

Construction

CPA-secure SKE
+
CMA-secure MAC $\Rightarrow$ AE $\left\{ \begin{array}{l} \text{CCA-secure} \\ \text{Unforgeable} \end{array} \right.$

- ① Encrypt-and-MAC
- ② MAC-then-Encrypt
- ③ Encrypt-then-MAC

Chosen-Ciphertext Attack (CCA) Security

$$k \leftarrow \text{Gen}(1^\lambda)$$

Alice



$$C_0 \leftarrow \text{Enc}_k(m_0)$$

$$C_1 \leftarrow \text{Enc}_k(m_1)$$

$$C_2 \leftarrow \text{Enc}_k(m_2)$$

⋮

$$m_0, m_1 \in \{0, 1\}^n$$

$$b \leftarrow \{0, 1\}$$

$$C \leftarrow \text{Enc}_k(m_b)$$

① Choose

② See

Choose

See

④ See

③ Choose

(PPT)

guess b'

$$\Pr[b' = b] \leq \frac{1}{2} + 2^{-\lambda}$$

Bob



$$m_0 := \text{Dec}_k(C_0)$$

$$m_1 := \text{Dec}_k(C_1)$$

$$m_2 := \text{Dec}_k(C_2)$$

⋮

$$m'_0 := \text{Dec}_k(C'_0)$$

$$m'_1 := \text{Dec}_k(C'_1)$$

$$m'_2 := \text{Dec}_k(C'_2)$$

⋮

$C_0, C_1, C_2, \dots$

C

$C'_0, C'_1, C'_2, \dots (\neq C)$

Unforgeability

$k \leftarrow \text{Gen}(1^\lambda)$

Alice



$c_0 \leftarrow \text{Enc}_k(m_0)$

$c_1 \leftarrow \text{Enc}_k(m_1)$

$c_2 \leftarrow \text{Enc}_k(m_2)$

$\vdots$

Bob



$m_0 := \text{Dec}_k(c_0)$

$m_1 := \text{Dec}_k(c_1)$

$m_2 := \text{Dec}_k(c_2)$

$\vdots$

$c_0, c_1, c_2, \dots$

① choose

② see



(PPT)

$\Rightarrow$ produce c^*

$\Pr[\text{Dec}_k(c^*) = m^*, m^* \neq m_i \forall i] \leq 2^{-\lambda}$

Encrypt-and-MAC

Given a CPA-secure SKE scheme $\Pi_1 = (\text{Gen}_1, \text{Enc}_1, \text{Dec}_1)$

& a strongly CMA-secure MAC scheme $\Pi_2 = (\text{Gen}_2, \text{Mac}_2, \text{Vrfy}_2)$

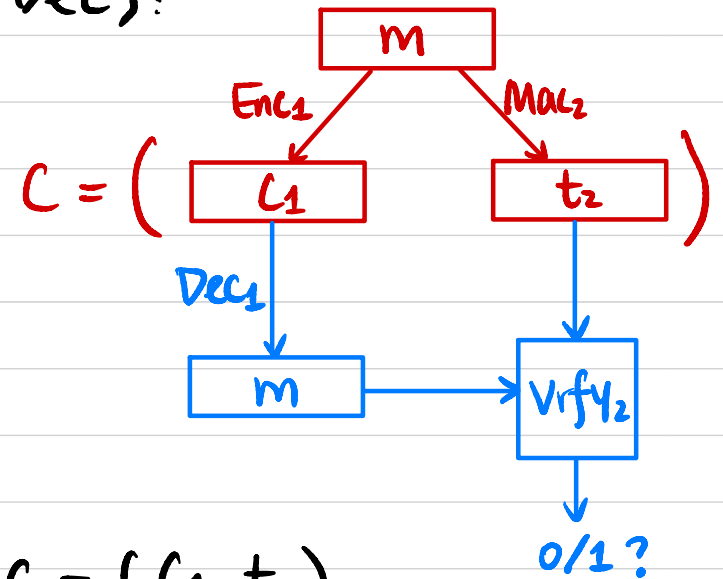
Construct an AE scheme $\Pi = (\text{Gen}, \text{Enc}, \text{Dec})$:

Gen(1^λ):

$k_1 \leftarrow \text{Gen}_1(1^\lambda)$

$k_2 \leftarrow \text{Gen}_2(1^\lambda)$

Output $k = (k_1, k_2)$



Enc_k(m):

$c_1 \leftarrow \text{Enc}_1(k_1, m)$

$t_2 \leftarrow \text{Mac}_2(k_2, m)$

Output $C = (c_1, t_2)$

Dec_k(C): $C = (c_1, t_2)$

$m := \text{Dec}_1(k_1, c_1)$

$b := \text{Vrfy}_2(k_2, (m, t_2))$

If $b=1$, output m

Otherwise output $\perp$

MAC-then-Encrypt

Gen(1^λ):

$$k_1 \leftarrow \text{Gen}_1(1^\lambda)$$

$$k_2 \leftarrow \text{Gen}_2(1^\lambda)$$

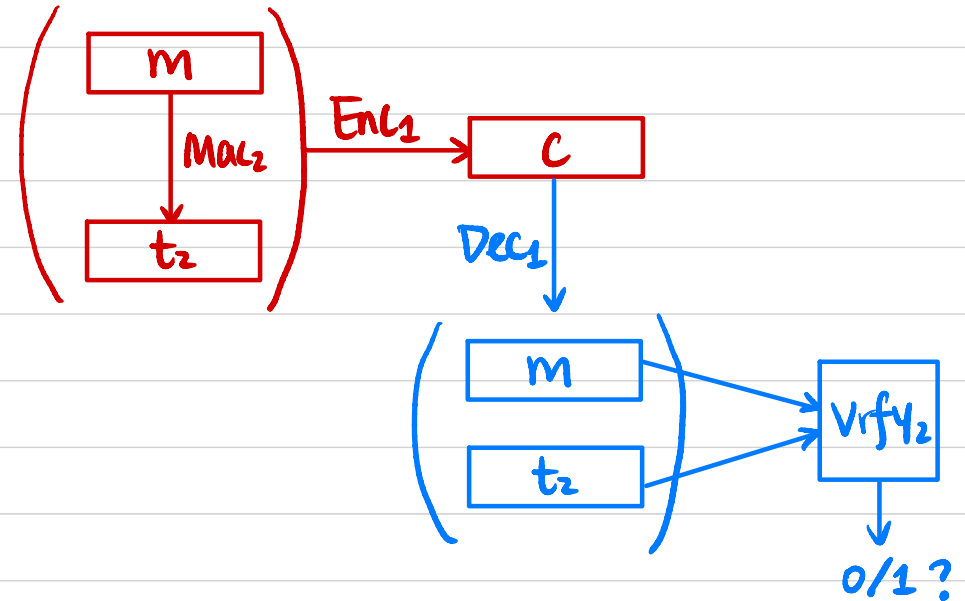
Output $k = (k_1, k_2)$

Enc $_k$ (m):

$$t_2 \leftarrow \text{Mac}_2(k_2, m)$$

$$c_1 \leftarrow \text{Enc}_1(k_1, m \parallel t_2)$$

output $C = c_1$



Dec $_k$ (C): $C = c_1$

$$m \parallel t_2 := \text{Dec}_1(k_1, c_1)$$

$$b := \text{Vrfy}_2(k_2, (m, t_2))$$

If $b=1$, output m

Otherwise output $\perp$

Encrypt-then-MAC

Gen(1^λ):

$$k_1 \leftarrow \text{Gen}_1(1^\lambda)$$

$$k_2 \leftarrow \text{Gen}_2(1^\lambda)$$

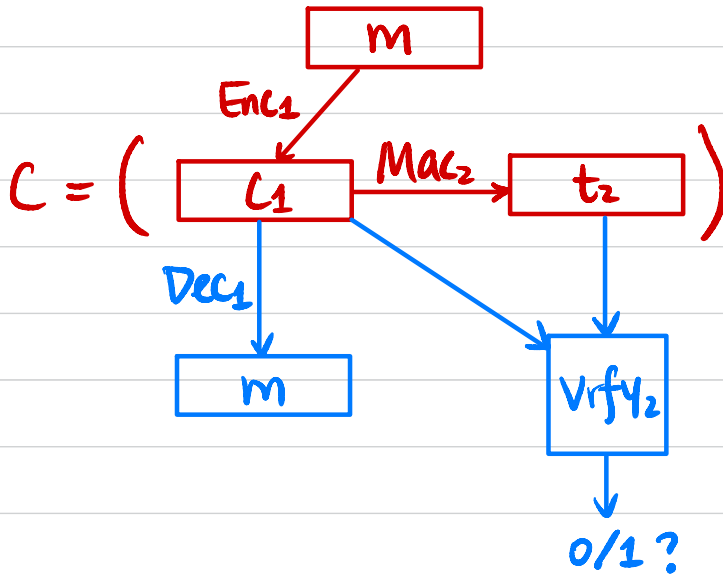
Output $k = (k_1, k_2)$

Enc $_k$ (m):

$$c_1 \leftarrow \text{Enc}_1(k_1, m)$$

$$t_2 \leftarrow \text{Mac}_2(k_2, c_1)$$

output $C = (c_1, t_2)$



Dec $_k$ (C): $C = (c_1, t_2)$

$$m := \text{Dec}_1(k_1, c_1)$$

$$b := \text{Vrfy}_2(k_2, (c_1, t_2))$$

If $b=1$, output m

Otherwise output $\perp$