

Agenda:

- Batching scheme we left off w/ last time
- Iteration-Level scheduling ("Continuous Batching")
- Different Inference Perf metrics: TTFT, TBT
- scheduling prioritization options (metric-dependent)
 - ↳ Decode vs. Prefill prioritizing
- chunked prefill

* Resources:

- Any scale BlogPost : <https://www.anyscale.com/blog/continuous-batching-llm-inference>
- ORCA OSDI Paper : <https://www.usenix.org/conference/osdi22/presentation/you>
- Sarathi Serve Paper <https://www.usenix.org/system/files/osdi24-agrawal.pdf>
- vLLM paper: <https://arxiv.org/pdf/2309.06180>

BATCHING SCHEME WE LEFT OFF WITH.

- Assumption: we are serving requests for a SINGLE model
- 7b model in fp16 takes up:
 $7 \times 10^9 \cdot 2 \approx 14 \text{ GB}$
- Approx. remaining of GPU memory (40 GB) can be used by KV cache
- LLM serving is quite memory bound (want to keep model weights loaded) - and serve

multiple (a BATCH of requests)

* note - we still don't want to waste

compute and recompute KV cache at each decode, so:

- when we serve a request, we want to make sure we can fit its entire KV

cache in GPU memory

↳ otherwise recompute would be too much

- maximum batch size we can fit is therefore determined by:

1. overhead per token in KV cache

2. max sequence length any request

can have

↳ (KV) $2 \times 2 \times \text{layers} \times d_{\text{model}}$

inex. from last class
24 1 2048

= .0002 GB 26 GB $\rightarrow$ mem. for KV cache = 130K

tokens that can fit = .0002 GB

max batchsize = $130K / \frac{(1024)}{\text{max seq length}} = 128$

"static" batching scheme we touched at:

1. wait until $\langle \text{batch size} \rangle$ requests are admitted to the system

2. For all batches, run prompt-processing step * to produce 1st token *

3. Proceed one "iteration" at a time, where we run 1 decode step across all our examples in the batch to predict the next token in each sequence.

↳ if at any point, a sequence ends,

don't do anything w/ that example

until all batch size * examples finish

* clarification: ^{w/o} paged attn memory management scheme from last time, * effective * batch size is lower (due to fragmentation, not all memory is available)

- say we have determined some batch size we know will fit, why is above scheme suboptimal?

* not all requests will use max # of tokens!!

<https://www.anyscale.com/blog/continuous-batching-llm-inference>

T_1	T_2	T_3	T_4	T_5	T_6	T_7	T_8
S_1	S_1	S_1	S_1				
S_2	S_2	S_2					
S_3	S_3	S_3					
S_4	S_4	S_4					

T_1	T_2	T_3	T_4	T_5	T_6	T_7	T_8
S_1	S_1	S_1	S_1	S_1	END		
S_2	S_2	S_2	S_2	S_2	S_2	S_2	END
S_3	S_3	S_3	S_3	END			
S_4	S_4	S_4	S_4	S_4	S_4	END	

white boxes = idle gpu cycles

* how often is this a problem (likelihood of white squares?)

↳ depends on whether input sequences can

together have **similar/same size**

↳ if all reqs produce 1 token (e.g., classification)

↳ if all prompts are same size

- in workloads like chatbots, there is a HIGH likelihood that sequences in same batch have diff. sizes

ITERATION-LEVEL SCHEDULING / CONTINUOUS BATCHING

- introduced by orca (OSDI 2022) - vllm implements this now

→ higher GPU utilization!

T_1	T_2	T_3	T_4	T_5	T_6	T_7	T_8
S_1	S_1	S_1	S_1				
S_2	S_2	S_2					
S_3	S_3	S_3					
S_4	S_4	S_4	S_4				

T_1	T_2	T_3	T_4	T_5	T_6	T_7	T_8
S_1	S_1	S_1	S_1	S_1	END	S_6	S_6
S_2	S_2	S_2	S_2	S_2	S_2	S_2	END
S_3	S_3	S_3	S_3	END	S_5	S_5	S_5
S_4	S_4	S_4	S_4	S_4	S_4	END	S_7

*handles late-arriving / early finish reqs more efficiently.

*key idea: let new requests start immediately if there is space

↳ at each iteration - check if a seq.

has ended and whether to admit a new request

INFERENCE PERF METRICS

* TTFT - time to 1st token → how long does prefill take

* TBT - time between tokens → how long does it take to generate each subsequent token?

- there are 2 parts to serving (token prediction tasks):

→ can be run in parallel (in one step)

1. Prefill: prompt processing to get 1st token

2. Decode: each subsequent iteration after

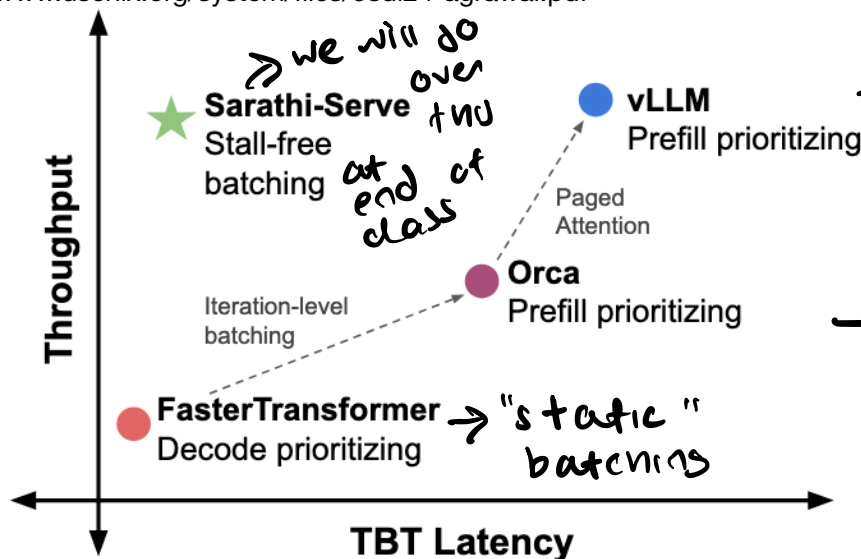
* prioritizing prefill will get you lower TTFT, ^{higher throughput} (fit more before)

* prioritizing ongoing decodes will get you lower

TBT, also will compromise throughput

* what does it mean to prioritize? Do these things come at a cost / is there a fundamental tradeoff?

<https://www.usenix.org/system/files/osdi24-agrawal.pdf>



w/ iteration-level scheduling - (continuous batching), we are:

- prioritizing prefill lets you fill in the empty space in batch, increasing throughput

(time for decodes)

*why does prefill prioritizing sacrifice

TBT / decode latency?

*let's look more closely at "static" batching

and "continuous batching" scheduling algorithms:

request-level

• **STATIC BATCHING**: optimizes TBT, but

wastes GPU utilization
(FasterTransformer)

<https://www.usenix.org/system/files/osdi24-agrawal.pdf>

```
1: Initialize current batch  $B \leftarrow \emptyset$ 
2: while True do
3:   if  $B = \emptyset$  then  $\rightarrow$  only when batch is empty,
                        do we admit a new request
4:      $R_{new} \leftarrow \text{get\_next\_request}()$ 
5:     while  $\text{can\_allocate\_request}(R_{new})$  do
6:        $B \leftarrow B + R_{new}$   $\hookrightarrow$  if it "fits" in batch
7:        $R_{new} \leftarrow \text{get\_next\_request}()$ 
8:      $\text{prefill}(B)$ 
9:   else  $\rightarrow$  if  $B$  has ANYTHING at all,
10:     $\text{decode}(B)$  finish on going decodes
11:     $B \leftarrow \text{filter\_finished\_requests}(B) \rightarrow$  check if  $B$  is empty
```

$\rightarrow$ **DECODE**
PRIORITIZING =
only admits
new req if
current
decodes
are
all
done
As we

slow before, lead
to low tput!

• what abt continuous batching?

```
1: Initialize current batch  $B \leftarrow \emptyset$ 
2: while True do  $\rightarrow$  at each iter
3:    $B_{new} \leftarrow \emptyset$   $\rightarrow$  initialize set of new req to
                        add
4:    $R_{new} \leftarrow \text{get\_next\_request}()$ 
5:   while  $\text{can\_allocate\_request}(R_{new})$  do
6:      $B_{new} \leftarrow B_{new} + R_{new}$   $\hookrightarrow$  checks there is
7:      $R_{new} \leftarrow \text{get\_next\_request}()$  enough space
                        for prefill of
                         $R_{new}$ 
8:   if  $B_{new} \neq \emptyset$  then
9:      $\text{prefill}(B_{new})$   $\rightarrow$  run all
10:     $B \leftarrow B + B_{new}$  NEW prefills
11:   else
12:     $\text{decode}(B) \rightarrow$  only run decodes once
13:     $B \leftarrow \text{filter\_finished\_requests}(B)$  we've finished new prefills
```

PREFILL PRIORITIZING:
at each
iteration,
run all
prefills we
can

↳ this makes TBT longer because we can have "generational stalls":

- before executing line 12 (decode),

we see if we have new prefix to run and run those first.

→ so there is a fundamental tradeoff b/w TBT latency and Throughput in existing schedulers.

- at each iteration, we have the choice to:

- ① run existing requests decode cycles -

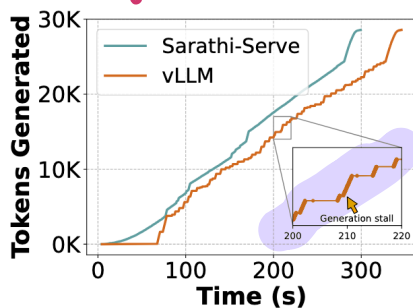
prioritizes TBT at expense of TBT

- ② run prefix - which will "stall ongoing decoder"

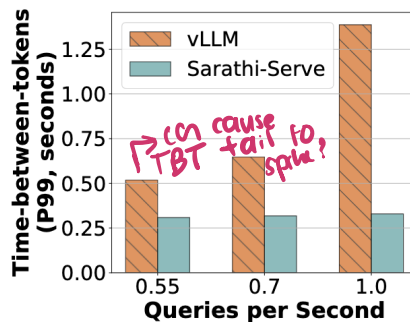
- what does this stalling look like?

- consider vllm in the graph below

↳ when decodes are interrupted, there will be a "stall" in new tokens generated



(a) Generation stall.



(b) High tail latency.

→ SARATHI SERVE / CHUNKED PREFILLS.

→ paper that analyzes existing inference

system and shows mismatch in tput / latency

→ proposes 2 techniques to solve:

1) chunked prefills

2) stall-free batching

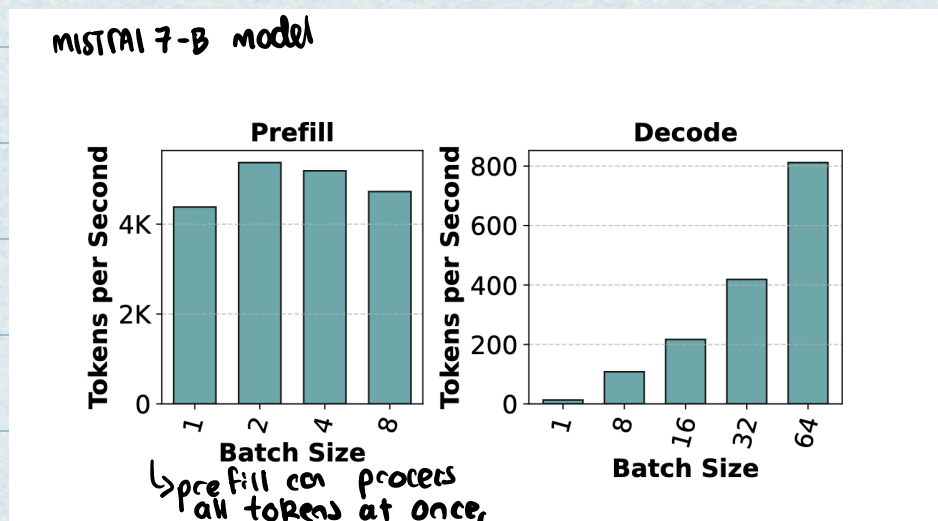
- we will try to build up these techniques

• 3 (related) observations:

→ OBSERVATION 1:

- batching boosts throughput for

decode, but not so much for prefill



↳ prefill can process all tokens at once, even batch size 1 is highly parallelizable

To get more tput

in decode, we need to

run more decodes.

Observation 2:

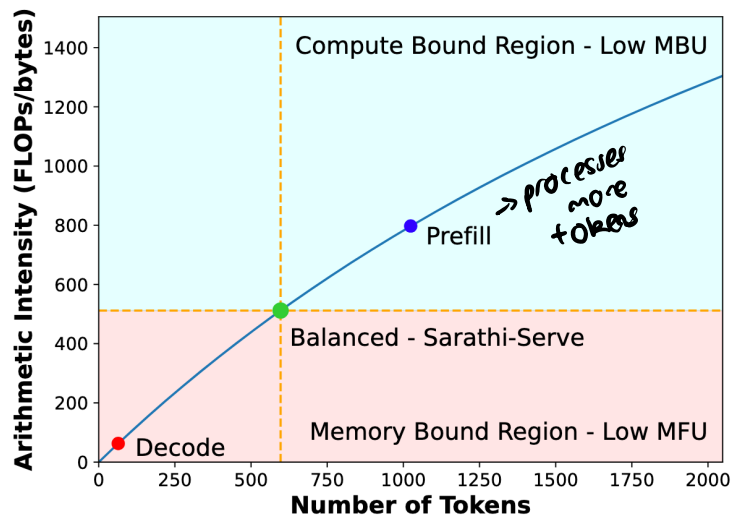
- Decode has low AI (memory bound),
low compute utilization

- Prefill has high AI (compute bound,
low memory bw utilization)

- neither at ridge point

these schemes use - more
Pipeline parallelism thoughts in paper

llm2-70B, running on 4A100s



- this graph JUST
focuses on linear layers -

so cost of fetching
layer weights vs. doing
computation on linear
layers

(another graph in paper
shows linear
ops as largest timespan)

- when we are in memory bound regions (e.g.,
decode) - we can do more compute for
same data fetched for free

↳ so *more tokens can be processed along
w/ decodes w/o running TBT for
ongoing decodes

- OBSERVATION 3:

- as we discussed before, existing system
choose at each iteration to run new

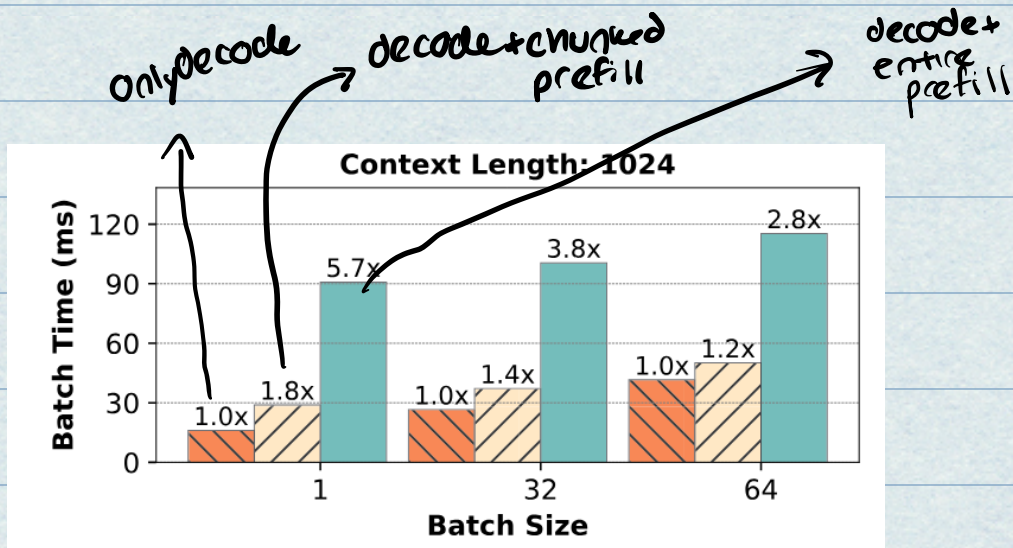
prefills (to inc. throughput) OR

finish existing requests

- running prefills can stall decodes

IDEA 1: chunked prefill

- when we have extra capacity from decodes, we'd like to do more work (decode has low AI, and is memory bound)
- if we try to add in a ENTIRE prefill - and we schedule a single iteration w/ existing decodes and a full prefill, latency will go up
- KEY IDEA: lets CHUNK the prefill:
 - ↳ and execute small blocks at a time



* we can execute a small portion of prefill along w/ decodes, to increase throughput w/o affecting latency that much

IDEA2: STALL-Free Batches

↳ w/ chunked prefills, how do we schedule each iteration?

```
1: Input:  $T_{max}$ , Application TBT SLO.
2: Initialize  $token\_budget$ ,  $\tau \leftarrow compute\_token\_budget(T_{max})$ 
3: Initialize  $batch\_num\_tokens$ ,  $n_t \leftarrow 0$ 
4: Initialize current batch  $B \leftarrow \emptyset$ 
5: while True do
6:   for  $R$  in  $B$  do
7:     if  $is\_prefill\_complete(R)$  then
8:        $n_t \leftarrow n_t + 1 \rightarrow 1^{st} \text{ decode}$ 
9:     for  $R$  in  $B$  do
10:      if not  $is\_prefill\_complete(R)$  then
11:         $c \leftarrow get\_next\_chunk\_size(R, \tau, n_t)$ 
12:         $n_t \leftarrow n_t + c \rightarrow \text{add next prefill for existing requests}$ 
13:       $R_{new} \leftarrow get\_next\_request()$ 
14:      while  $can\_allocate\_request(R_{new}) \wedge n_t < \tau$  do
15:         $c \leftarrow get\_next\_chunk\_size(R_{new}, \tau, n_t)$ 
16:        if  $c > 0$  then
17:           $n_t \leftarrow n_t + c$ 
18:           $B \leftarrow R_{new}$ 
19:        else
20:          break
21:
22:       $process\_hybrid\_batch(B)$ 
23:       $B \leftarrow filter\_finished\_requests(B)$ 
24:       $n_t \leftarrow 0$ 
```

paper has more details on calculations tho, but we can think of as additional capacity until we become compute bound

→ only add new prefill chunks after
(a) existing decodes
(b) ongoing chunked prefills

→ prioritize ongoing decodes: Always schedule these 1st to run

→ THEN schedule next chunk of existing batches

→ THEN only admit new requests

*process hybrid batch.

- can execute batch of decodes w/
chunked prefills at same time
- previous systems were only executing
a batch of decodes OR batch of
prefills
- paper has ablations on hybrid batch
technique

* some related themes of work:

- because prefill and decode behave so differently,
why not put them on dif. GPUs? scheduling
will be easier
 - ↳ can send KV cache in between
- lets cache KV cache in CPU memory
(when it needs to be reused) - e.g. documents
and transfer it
 - ↳ but compress it in a smart
way so its not so large to send