

# Set Equality & Quantification

Michael L. Littman

CS 22 2020

January 29, 2020

# Overview

Proof by Cases (1.7)

Predicate Formulas (3.6)

## Fact about groups of people

Any two people have either met or not.

Given a group of people  $G$ , if all pairs of people in  $G$  have met, we'll call it a *club*. If no two people in  $G$  have met, we'll call them *strangers*.

**Theorem.** Every collection of 6 people includes a club of 3 people or a group of 3 strangers.

Does that seem true? Try some examples on the board.

## Proof (Part 1)

The proof is by case analysis. Let  $x$  denote one of the six people. Let  $R = G - \{x\}$  be the rest. There are two cases:

1. Among  $R$ , at least 3 have met  $x$ .
2. Among  $R$ , at least 3 have *not* met  $x$ .

At least one of these cases must hold. Since  $|R|$  is odd, either more than half in  $R$  know  $x$  or less than half in  $R$  know  $x$  (and therefore more than half do not know  $x$ ).

Case 1: At least 3 have met  $x$ . Let  $J \subseteq R$  be those individuals.

Two subcases:

- 1.1 No pair in  $J$  have met each other. So,  $J$  is a group of at least 3 strangers and the theorem holds in this subcase.
- 1.2 Some pair in  $J$  have met each other. That pair and  $x$  are a club of 3 people and the theorem holds in this subcase, too.

That covers Case 1!

## Proof (Part 2)

Case 2: At least 3 have not met  $x$ . Let  $J \subseteq R$  be those individuals.  
Two subcases:

- 2.1 Every pair in  $J$  have met each other. So,  $R$  is a club of at least size 3 and the theorem holds in this subcase.
- 2.2 Some pair in  $J$  haven't met each other. That pair and  $x$  are a group of strangers of 3 people and the theorem holds in this subcase, too.

That covers Case 2! It's kind of the inverse-video version of Case 1.

Since we showed that only these two cases can occur and the theorem holds in both, the theorem *always* holds.

## Quantifiers, Revisited

**Always True** (universal quantification)

$$\forall x \in \mathbb{R}, x^2 + 1 \geq 0.$$

- ▶ For all  $x \in D$ ,  $P(x)$  is true.
- ▶  $P(x)$  is true for every  $x$  in the set  $D$ .

**Sometimes True** (existential quantification)

$$\exists x \in \mathbb{Z}, x \text{ is even and } x \text{ is prime.}$$

- ▶ There is an  $x \in D$  such that  $P(x)$  is true.
- ▶  $P(x)$  is true for some  $x$  in the set  $D$ .
- ▶  $P(x)$  is true for at least one  $x \in D$ .

## Mixing quantifiers

**Theorem** (sparse squares): There's a perfect square arbitrarily far from its closest perfect square.

Clear? Maybe a tad vague. True? How say in math?

$\forall d \in \mathbb{N}, \exists i \in \mathbb{N}, \forall j \in \mathbb{N}, i$  is a perfect square AND  $|i - j| \leq d$   
IMPLIES  $j$  is NOT a perfect square.

The expressions nest inside each other. The order matters.

You can think of it like a little game. I'm claiming that you can pick any  $d$  you want. I'll then pick an  $i$  that's a perfect square AND no matter what  $j$  you pick that is within  $d$  values of  $i$ ,  $j$  won't be a perfect square.

So, what's my winning strategy?

## Any ambiguity is too many

“If you can juggle any object, you’ve got a talent.”

1. If  $\exists o$ , you can juggle  $o$ , then you’ve got a talent.
2. If  $\forall o$ , you can juggle  $o$ , then you’ve got a talent.

“...statistics show that, in New York, a man is mugged every 11 seconds. I would now like you to meet that man. His name is Jesse Donnally.”

1.  $\forall t, \exists m$ ,  $m$  mugged at time  $t$
2.  $\exists m, \forall t$ ,  $m$  mugged at time  $t$



## From my files

Addressing a group: “Send me all of your papers.”

- ▶  $\forall x \text{ in group, } \forall \text{ papers } p, x \text{ wrote } p \text{ IMPLIES send}(x)$
- ▶  $\forall \text{ papers } p, (\forall x \text{ in group, } x \text{ wrote } p) \text{ IMPLIES send}(x)$

About a medical side effect: “Everything tastes the same”

- ▶  $\forall x, \forall y, \text{ taste}(x, \text{now}) = \text{taste}(y, \text{now})$
- ▶  $\forall x, \text{ taste}(x, \text{now}) = \text{taste}(x, \text{then})$

“The whole article is not available.”

- ▶  $\text{not } \forall \text{ article part } x, x \text{ is available}$
- ▶  $\forall \text{ article part } x, x \text{ is not available}$

## DeMorgan returns: Negating quantifiers

These two statements are equivalent:

- ▶ Not everyone likes chocolate.
- ▶ There's someone who doesn't like chocolate.

not  $\forall x, P(x)$  is equivalent to  $\exists x, \text{not } P(x)$ .

## Assertion about predicates

$\exists x, \forall y, P(x, y)$  IMPLIES  $\forall y, \exists x, P(x, y)$ .

If  $\exists x, \forall y, P(x, y)$ , there must be some specific  $x^*$  such that  $\forall y, P(x^*, y)$ . As a result,  $\forall y, \exists x, P(x, y)$  because we can always choose  $x^*$  to be the selected  $x$ .

On the other hand,

$\forall y, \exists x, P(x, y)$  IMPLIES  $\exists x, \forall y, P(x, y)$

is not true.

	$y_1$	$y_2$	$y_3$
$x_1$	T	F	F
$x_2$	T	T	T
$x_3$	F	T	F

	$y_1$	$y_2$	$y_3$
$x_1$	T	F	F
$x_2$	T	F	T
$x_3$	F	T	F