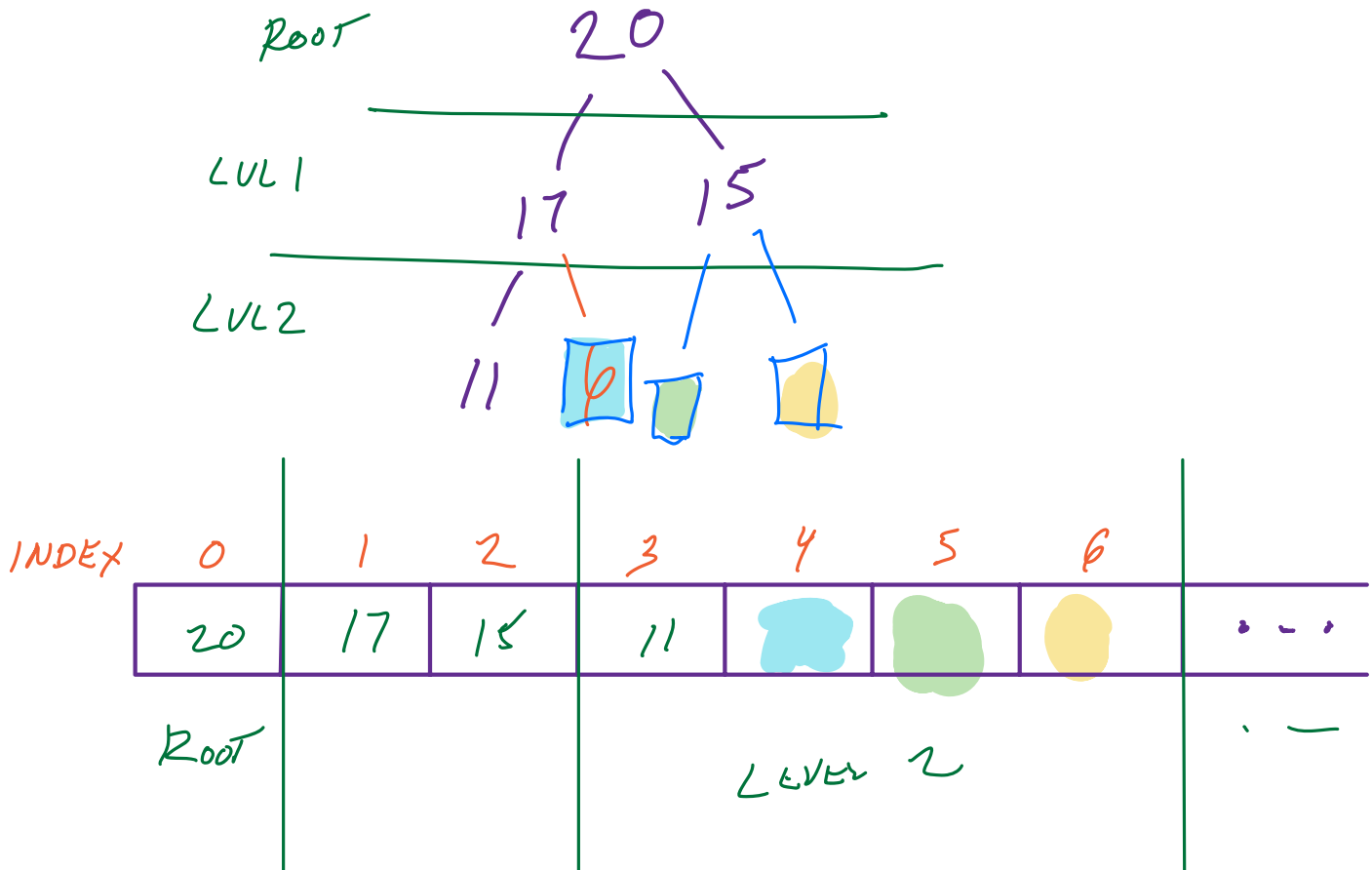


From last time: how to implement a heap?

Want a data structure that:

- Easy to find open slot => insert
- Easy to navigate up and down tree => help with inserting/removing elements

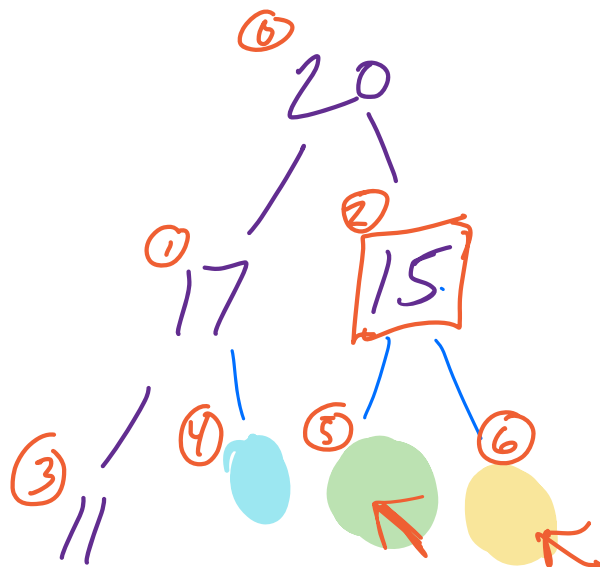
SOLUTION: EMBED TREE INTO AN ARRAY



```
tree = [20, 17, 15, 11]
tree.append(6)
[20, 17, 15, 11, 6]
```

Can implement insert by adding to the end of the array  
=> always fills in the next available slot!

## How does this array structure make a tree?



Find next open slot: just look for first empty element

For some node at index  $i$ :

- left( $i$ ) is at  $(i * 2) + 1$
- right( $i$ ) is at  $(i * 2) + 2$
- parent( $i$ ) is at  $\text{floor}((i - 1)/2)$

Ex.  $i = 2$  holds 15

$$\text{LEFT} = 2 * 2 + 1 = 5$$

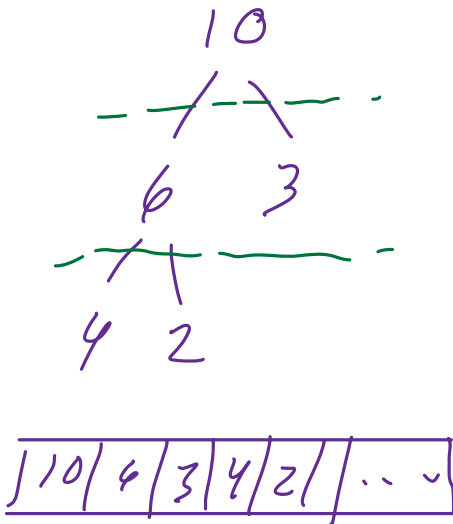
$$\text{RIGHT} = 2 * 2 + 2 = 6$$

$$\text{PARENT} = \left\lfloor \frac{2-1}{2} \right\rfloor = \left\lfloor \frac{1}{2} \right\rfloor = 0$$

=> Can leverage array structure to move between elements in the tree, all in constant time!

INDEX	0	1	2	3	4	5	6	
	20	17	15	11				
Root	LEVEL 1		LEVEL 2					

## Recap: Embedding trees/heaps in Arrays



For some node at index  $i$ :

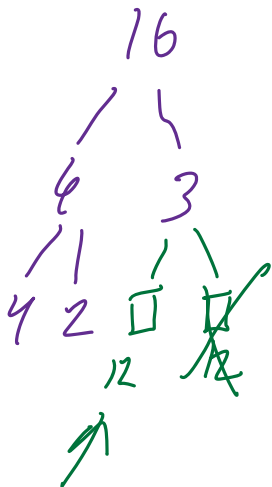
- $\text{left}(i)$  is at  $(i * 2) + 1$
- $\text{right}(i)$  is at  $(i * 2) + 2$
- $\text{parent}(i)$  is at  $\text{floor}((i - 1)/2)$

How to insert?

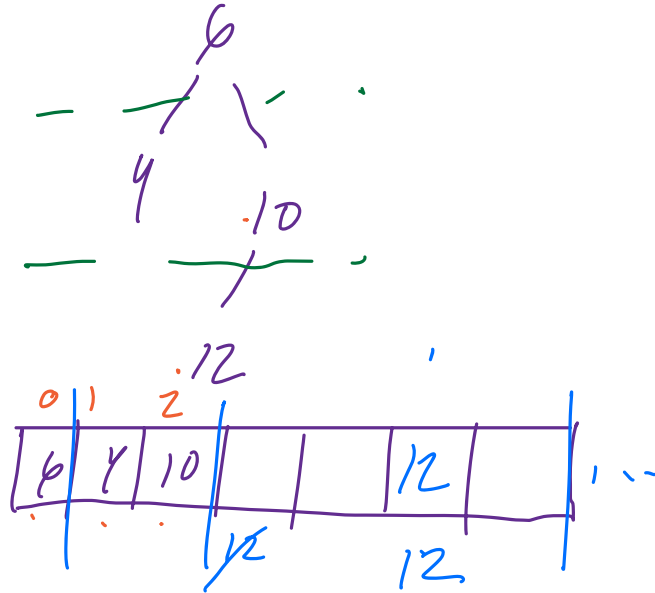
In the code version: always add new elements to the next available space

=> Keeps balanced

=> Always know how to find next slot (just the end of the array)



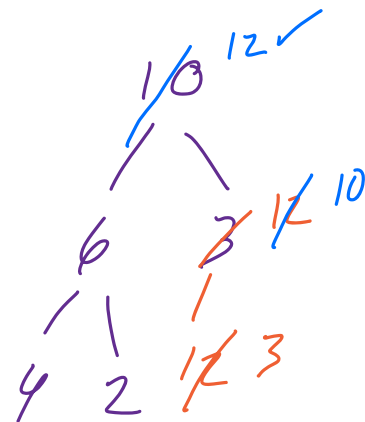
WHAT ARE THE ARRAYS  
FOR THESE TREES?



Why not keep stuff condensed? Need to respect parent/child formulas to match the structure of the tree

=> Without spaces, parent, left/right formulas will break

Ex. INSERT 12

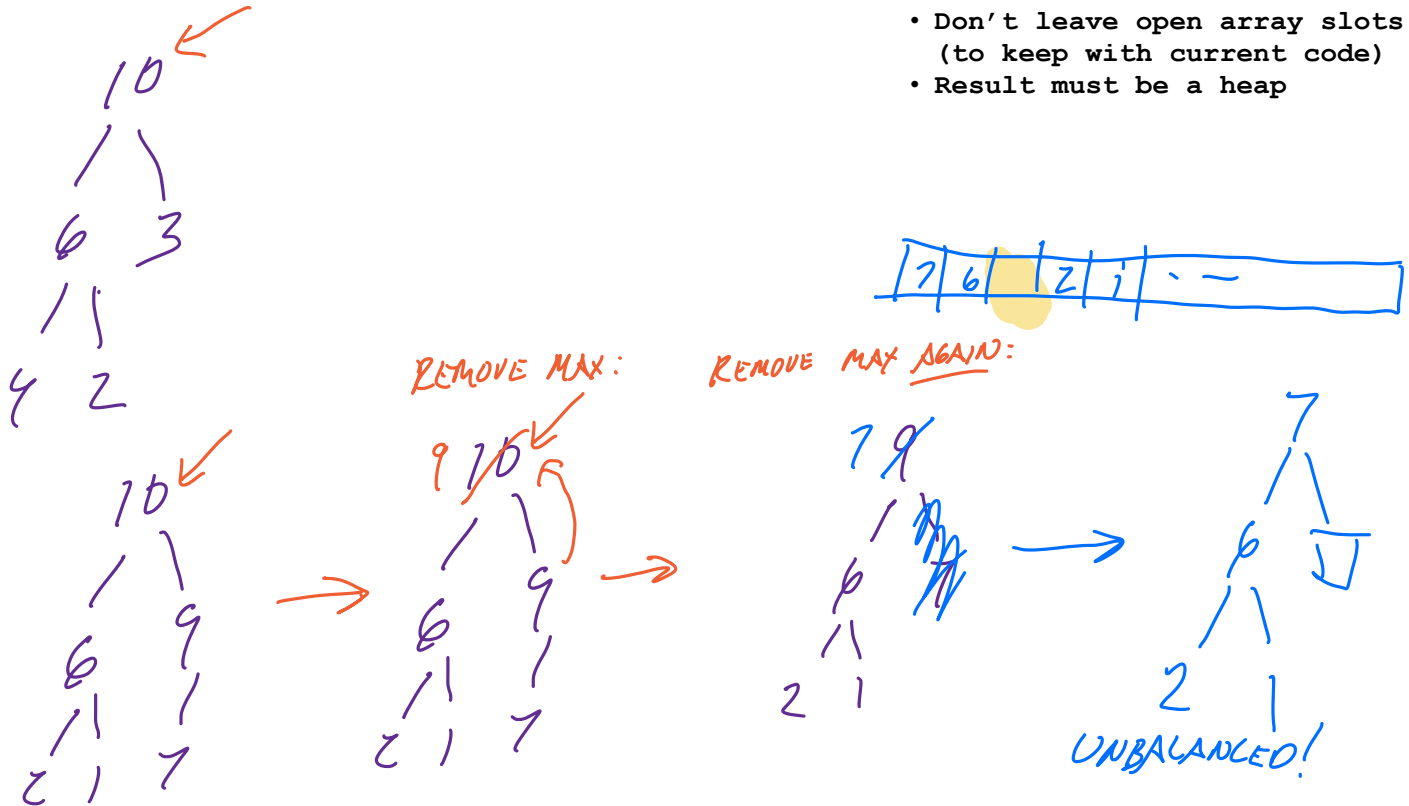


How would we implement delete\_max?

Let's try something similar to insert, where we swap nodes up after removing an element:

#### Requirements

- Maintain balance
- Don't leave open array slots (to keep with current code)
- Result must be a heap

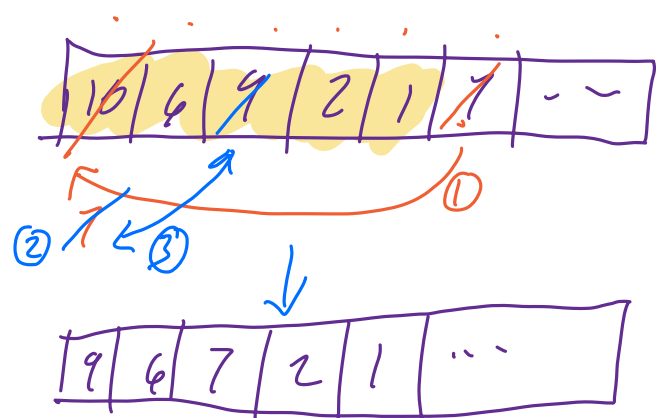


Problem: If use the same remove + swap from here, can end up with unbalanced tree  
=> Need to do delete without creating gaps  
=> If we can just swap values within existing cells in use, can avoid gaps  
=> Leverages array structure!

We didn't implement remove\_max, but here's the intuition:

- If we're removing an element, the resulting array must be one smaller than before
- Removing the max element creates a hole => swap the last element (ie, the one that would get "abandoned" if we were to shrink the array) to the top (in the hole left by the max)
- Swap this new top element down until the result is a heap

This idea leverages the same principles as insert, but involves swapping in a different direction. Similar to how we knew where to add a new element when inserting into the heap, we leverage the array to know where to find a new element to start swapping down (ie, last one)



# RECAP: NOW TO THINK ABOUT ARRAYS

Essential: have items in predictable, and computable, locations in memory

=> "where is the i'th element"

Hash Map  
Hash Set

Use predictable location to get from hash value to some specific index

=> Not all indices are used  
=> Positions correspond to "array slots"

Array List

Use predictable location for get(i)

=> Items in consecutive locations in memory

↑  
USUALLY  
THINK ABOUT  
SHIFTING ELEMENTS  
TO MAINTAIN THIS

Heap (Tree)

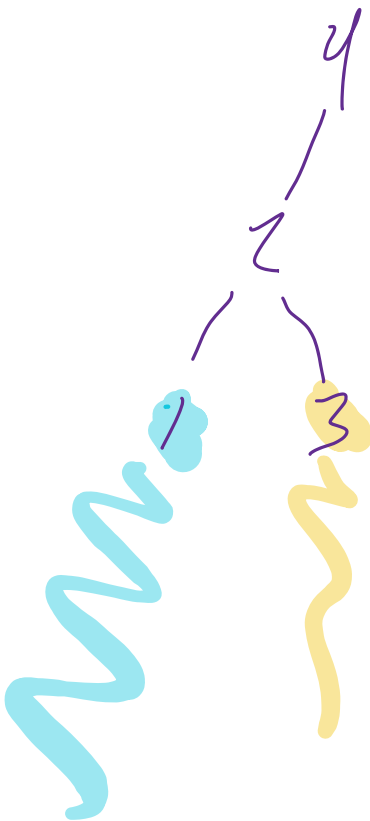
Use predictable location to navigate between parent/child nodes

=> Positions correspond to where we are in tree

Important to think about

- Ways different data structures can be used
- How the underlying data structure (arrays, in this case) can matter for different applications

Do all binary trees belong in arrays? What about BSTs?



What's questions do you need to ask to decide if this is a good idea?



Here's a way this tree could be balanced. Try it: think about how the array structure would need to change if we made this change.

What if (1) and (3) had lots of child nodes? With an array, we'd need to do a lot of shifting nodes around to make this work! For cases like this, it makes more sense to use a node-based representation (ie, the traditional recursive data structure approach we've seen before), where we just need to alter the parent/child references.

Overall: think about how are going to use the data structure  
=> What operations do you need to perform (at a high level)  
=> For BSTs: need to keep the BST ordered and balanced

=> How does that translate into operations on the data structure  
=> If we used an array, we would need to shift a lot of elements around to keep the BST balanced!

```

import math

"""implementation of a max heap"""
class Heap:
    def __init__(self):
        self.data = []
        self.size = 0

    def __str__(self):
        """string representation is the underlying list"""
        return str(self.data)

    def parent_index(self, of_index):
        """compute parent index of given index. Assumes of_index > 0"""
        return math.floor((of_index - 1) / 2)

    def swap(self, index1, index2):
        """swaps values in index1 and index2 within self.data"""
        tmp = self.data[index1]
        self.data[index1] = self.data[index2]
        self.data[index2] = tmp

    def insert(self, new_elt):
        """insert element into the heap"""
        self.data.append(new_elt)
        self.sift_up(self.size)
        self.size += 1

    def sift_up(self, from_index):
        """swap element in from_index up heap until it is in the right place"""
        if from_index > 0:
            parent = self.parent_index(from_index)
            if self.data[from_index] > self.data[parent]:
                self.swap(parent, from_index)
                self.sift_up(parent)

    def sift_up_while(self, from_index):
        """a while-loop based version of sift_up"""
        if from_index > 0:
            curr_index = from_index
            parent = self.parent_index(curr_index)
            while (curr_index > 0) and \
                (self.data[curr_index] > self.data[parent]):
                self.swap(parent, curr_index)
                curr_index = parent
                parent = self.parent_index(curr_index)
            # Note: this version repeats last two lines, unlike the recursive one

```