Design Problem: A hospital emergency room needs to manage information about the patients who need to be seen. Patients are treated in order of urgency (most urgent first). As each person comes in, a record is created with their name, urgency level (a number, in which higher numbers have higher priority), and a brief description of their injury.

Think about the kinds of operations that we have studied on sequences of data this semester:

- add J WANT TO DO THESE QUICKLP
  remove J GET QUICKLY, <u>BY URGENCY</u>
  contains MAYBE LESS IMPORTANT

Which of these seem most important when managing the order in which to treat patients? Which operations need to be the most efficient?

### Priority queue (PQ): 3 fundamental operations

- 1. Insert a new item
- 2. Remove the maximum priority item
- 3. Get maximum priority item, without removing (peek)

#### => These are the operations we want to prioritize

How do we build such a data structure from scratch? Consider the data structures that we have studied so far: which seem to be suitable here?

- Linked List
- Array List
- Tree/Binary Search Tree
- Hashmap/Dictionary
- HashSet/set
- Class of our own design (e.g., graph)

Write the type of your proposed data structure for managing hospital triage data (in terms of types)

NEXT PAGE ...-



How do we build one from scratch?

Some data structure: arrays, lists, trees, hash maps/dictionary, hash sets/ set, graphs

How might you use these? Which ones might be best?

# HashSets/set

=> Not really a way to specify priority

# HashMap/dict

| Linked list       | keep a sorted list  |
|-------------------|---|
| <u>Array list</u> | - get-max: O(1) (pick first element)  |
|                   | - remove-max: O(1) for linked list, for array list would need to shift elements |
|                   | to keep sorted order => O(N)  |
|                   | - insert: O(N) to find position in sorted list                                  |

# <u>Trees</u>

We know one way to do an ordered representation with trees...

=> BST (Binary search tree): for any node, every smaller node is on the left, any larger node is on the right

What if we used this as a priority queue?



insert, get-max : O(logN) if balanced O(N) if unbalanced

BASICALLY A LIST 10

What if we relax the rules a bit?

=> For a priority queue, we don't need a total order like a BST. What if we just keep the max item at the top??

Heap (binary max heap): a binary tree (NOT a BST) with two constraints:

- max item is at the root
- left and right subtrees are also heaps

DATA: 2, 4, 8, 9, 10, 12



Note: can have different valid representations for the same heap (may be more or less-balanced... more on this later)

ALSO A NEAP

DATA: [1, 2, 5, 7, 8] NEAP EQUIVALEN NOTA HEAP! THIS SUBTREE SHOULD NEAP NAVE 5 AS MAX (NOT BALANCED, THOUGN)

Checking in on our priority queue goals: what can we infer about the runtime of using a heap for a PQ?

#### Priority queue (PQ): needs 3 operations

- 1. Insert a new item => ???
- 2. Remove max-priority item => ???
- 3. Get max-priority item (without removing) => O(1) => can just look at top of heap!

HOLE " FROM REMOVING 12 What about add and remove? EX, REMOVE 12 RESULT 10 10 NEXT CANDIDATES Removing an element creates a "hole", can reorder subtree to fill it To reorder, we only need to consider one "branch" of the heap => If heap is balanced, this takes O(logN)

WHAT IF WE WANT YO INSERT !!!

Strategy: add to bottom, reorder until we have a heap again

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Again, we only need to reorder one "branch" of the heap => O(logN)



Priority queue (PQ): needs 3 operations

#### if heap is balanced:

- 1. Insert a new item => O(logN)
- 2. Remove max-priority item => O(logN)
- 3. Get max-priority item (without removing)
   => 0(1)

If heap is unbalanced, the insert/delete steps are harder:

- insert: O(N)
- remove\_max: O(N)
- get\_max: 0(1)

<u>Open questions (for next time):</u>

- How to find an empty spot to insert?
- How to keep the heap balanced to ensure logN runtime?
- What does "balanced" even mean, anyway?