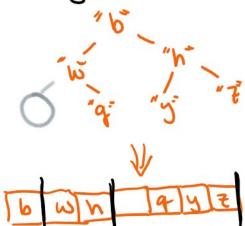


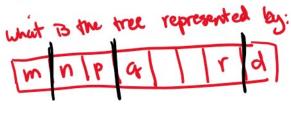
- a balanced heap reduces runtime want height to be log(# elts)
-what if we inserted systematically?
"row by-row", "left-to-right"

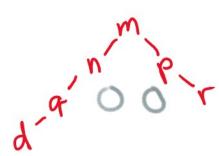
Challenge: if we use our BinTree class, how do we find the next empty spot to insert in? The possible solution has a downside (note that this only considers inserting into an empty spot, NOT swapping upwards): def insert anywhere(self, new elt): ''' inserts into an empty spot in the tree ''' if (not self.left): # found empty spot in left subtree self.left = BinTree(new elt) # still need to swap up elif (not self.right): # found empty spot in right subtree self.right = BinTree(new_elt) # still need to swap up else: self.left.insert_anywhere(new_elt) # find a spot in the left subtree # could have also done: self.right.insert anywhere(new elt) This "biases" the tree to the left (aka leads to an imbalanced tree). We could try to compare the heights of the left and right subtrees and insert into the smaller one, but that leads to complicated code.

Thed to do w our ... rewrite class ?

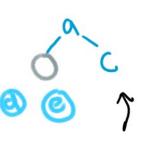
array representation of a lineary mee







ablolde



not all arrays can represent ratid trust



We can get an element's parent and children in constant time:

inter: 0 1 2 3 4 5 6.

If we want to enforce heap structure in the way we said, this array representation will have no "holes" as long as we always insert in the first available slot of the array and "sift up"! This means this representation leads to a balanced and easy-to-insert-into heap!!

lest child of ett. @ ind. i is @ ind 24it1
right child is @ ind 22it2
parent@ floor (21)