

The Swine Before Perl

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Brown University and PLT



Why We're Here

Any sufficiently complicated C or Fortran program contains an ad hoc, informally-specified, bug-ridden, slow implementation of half of Common Lisp.

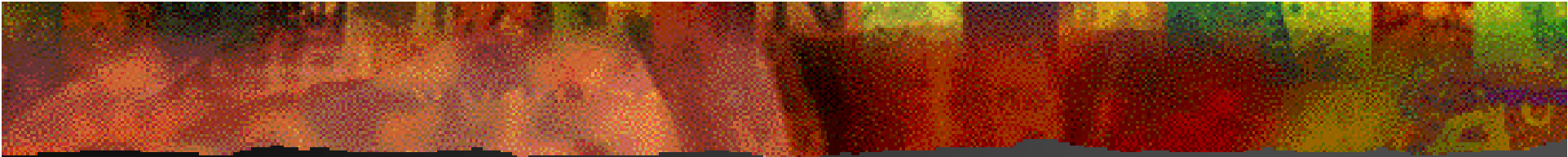
—Phil Greenspun's
Tenth Law of Programming

Our Corollary

Any sufficiently useful C or Fortran program
needs to be “sufficiently complicated”

We’re here to provide the heavy lifting

I’m here to motivate how Scheme does this



The Swine: PLT Scheme

“Whatever: can you
program in it?”

Books

- *Teach Yourself Scheme in Fixnum Days*
 - *How to Use Scheme*
 - Extensive manuals and user support
 - It's all **free of cost**
-

Software #1: MzScheme

- Quick, textual shell: ideal for scripting
 - Rich module system
 - Rich mixin-based object system
 - Embeddable in applications
 - Garbage collection across C/C++/Scheme
 - Libraries for lots of 3- and 4-letter acronyms (XML, CGI, ODBC, COM, ...)
-

 MzScheme

Welcome to MzScheme version 103, Copyright (c) 1995-2000 PLT (Matthew Flatt)

> (require-library "function.ss")

> (require-library "url.ss" "net")

>

(foldl + 0 (map file-size (directory-list)))

3887958

>

(call/input-url (string->url "http://www.cs.brown.edu/~sk/test.html")
get-pure-port display-pure-port)

<html>

<head><title>Test</title></head>

<body><p>Hello!</p></body>

</html>

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3887958

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(call-with-input-uri (string->uri "http://www.cs.brown.edu/~sk/test.html")
get-pure-port display-pure-port)

<html>

<head><title>Test</title></head>

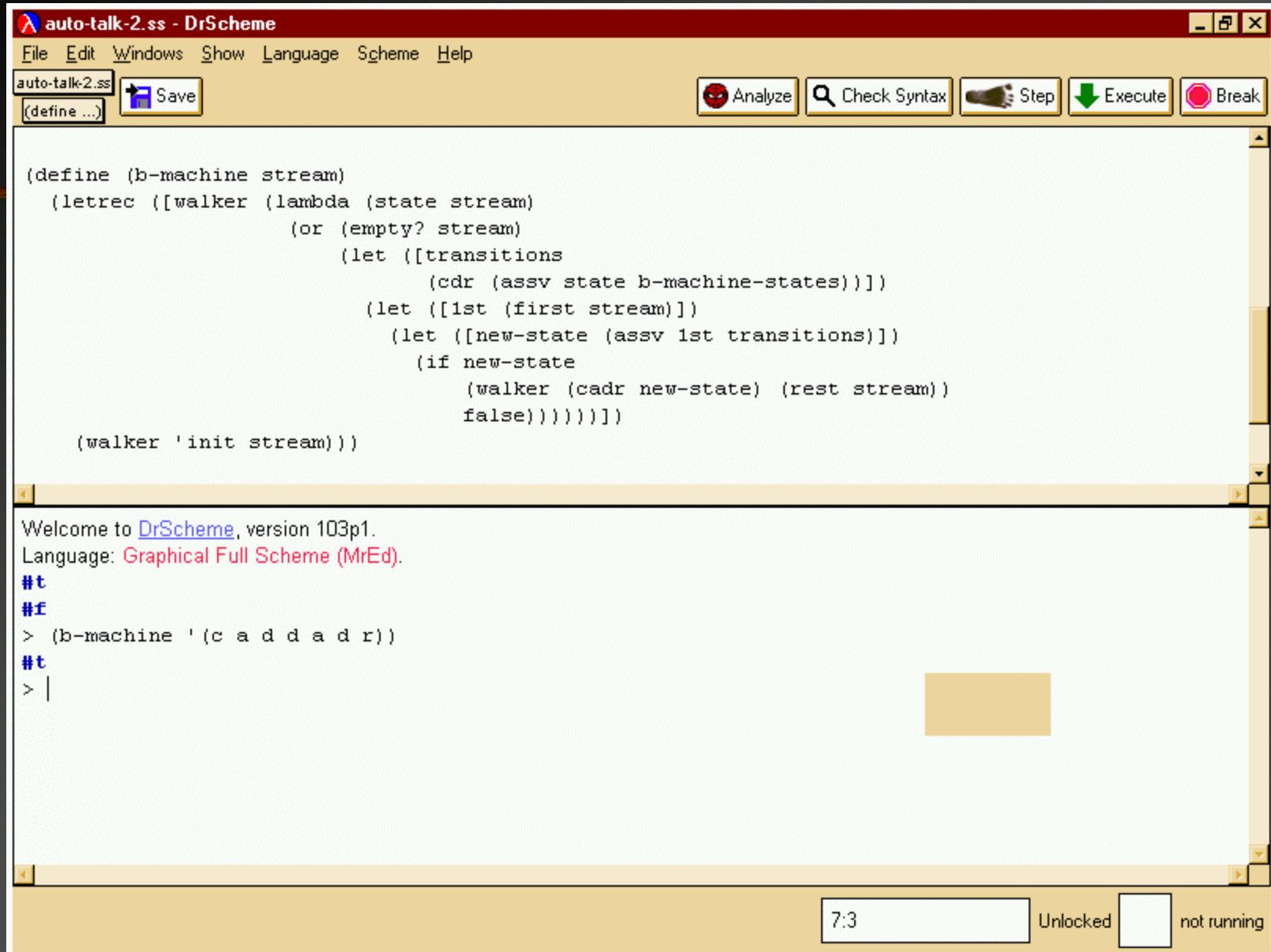
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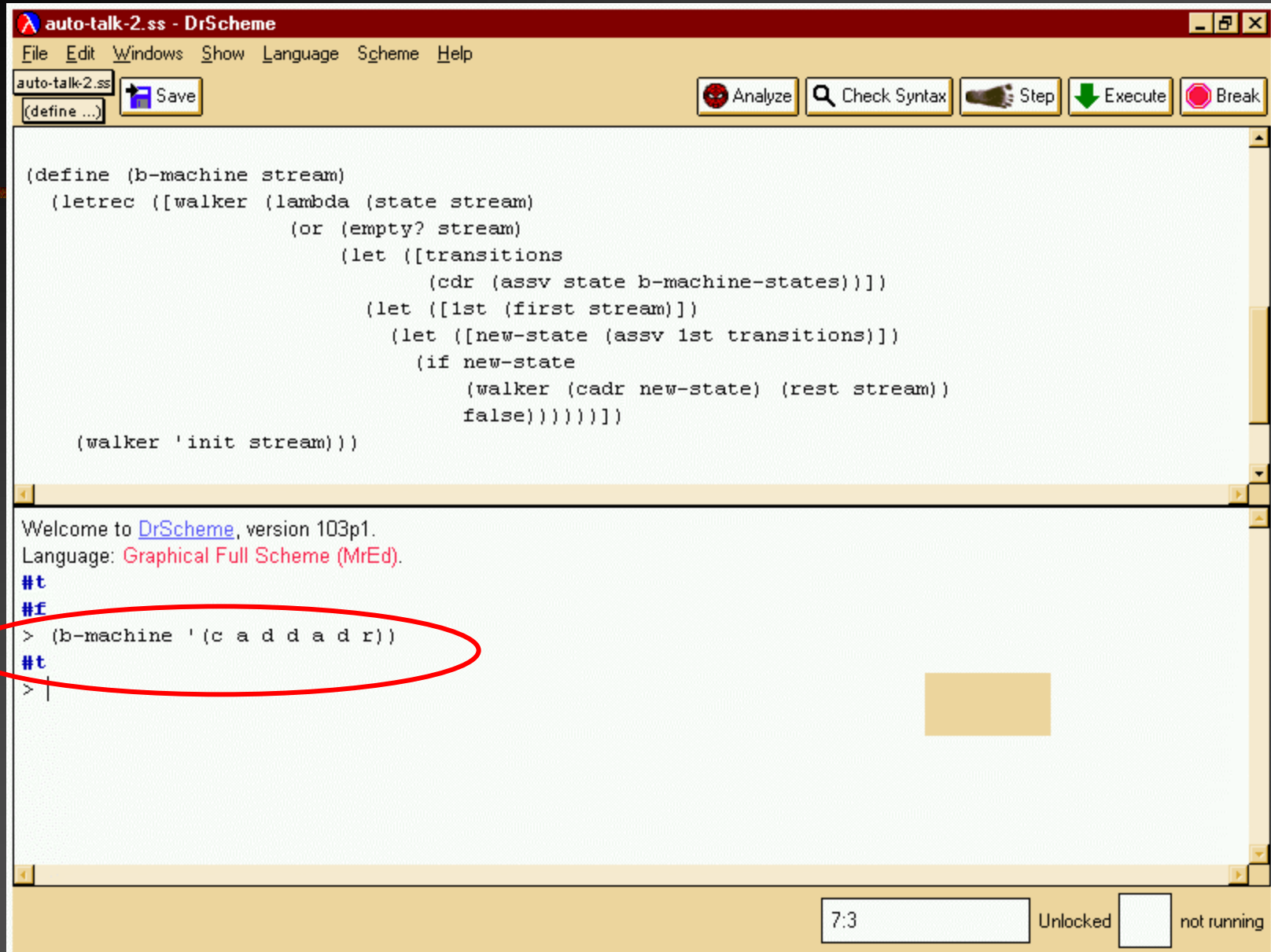
</html>

>

Software #2: DrScheme

- Fully-graphical programming environment
 - Full portability: Unix, Mac and Windows
 - Special support for beginning Schemers
-





Software #3: Type Inference

- Two PhD theses and counting
 - Graphically displays inferred types
 - Let's see how it works ...
-

auto-talk-2.ss - MrSpidey

File Edit Windows Actions Show Clear Filter Help

```
(define b-machine-states
  '((init (c loop))
    (loop (a loop)
          (d loop)
          (r end))
    (end (r end))))

(define (b-machine stream)
  (letrec ([walker (lambda (state stream)
                    (or (empty? stream)
                        (let ([transitions
                              (cdr (assv state b-machine-states))])
                          (let ([1st (first (car stream))])
                            (let ([new-state (assv 1st transitions)])
                              (if new-state
                                  (walker (cadr new-state) (rest stream))
                                  false)))))))]))

  (walker 'init stream)))

(b-machine '(c a d a d r))
(b-machine '(c a d a r d))
```

Welcome to [MrSpidey](#), version 103p1.

CHECKS:

[car](#) check in file "auto-talk-2.ss": line 1, column 15

[cdr](#) check in file "auto-talk-2.ss": line 18, column 34

[car](#) check in file "auto-talk-2.ss": line 19, column 47

Unlocked

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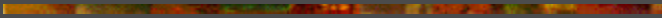
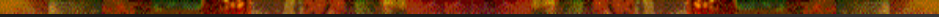
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Unlocked

Software #4: Web Server

- Dynamic content generation is a breeze
 - HTML/XML transformation is a breeze
 - Trivially build use-once security policies
 - For dynamic content, 8x speed of Apache
-

The Gems



The Gems

- Closures
 - Continuations
-

The Gems

- Closures
- Continuations

The Crown Jewels

The Gems

- Closures
- Continuations

The Crown Jewels

- That stupid parenthetical syntax
 - Macros
 - Tail calls
-

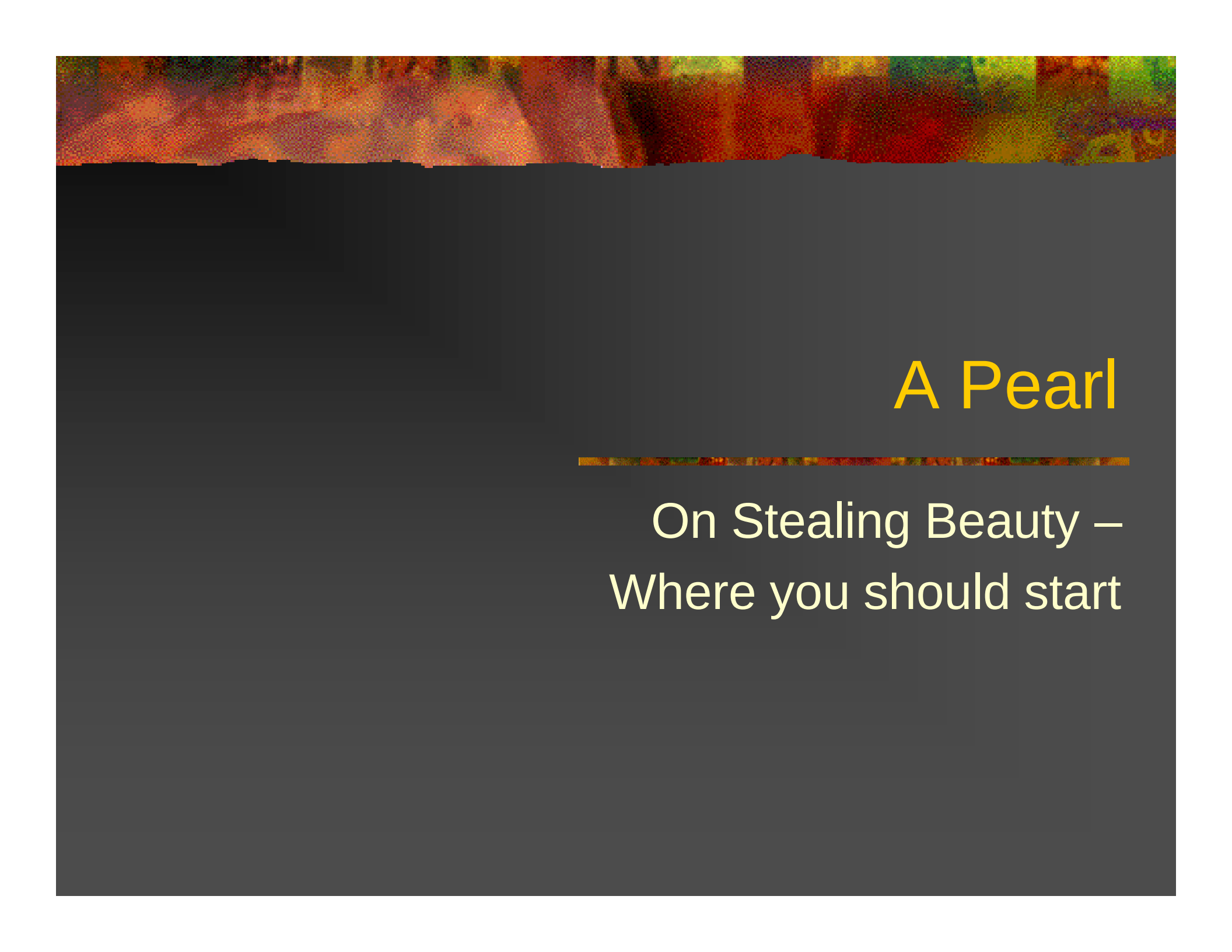
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- Closures
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The Crown Jewels

- That stupid parenthetical syntax
- Macros
- Tail calls

The
Lesser
Gems



A Pearl

On Stealing Beauty –
Where you should start

Problem

Pattern-matcher for streams:

- Consumes a stream of input tokens
 - Must be easy to write and read
 - Must be fairly fast
 - Must integrate well into rest of code
-

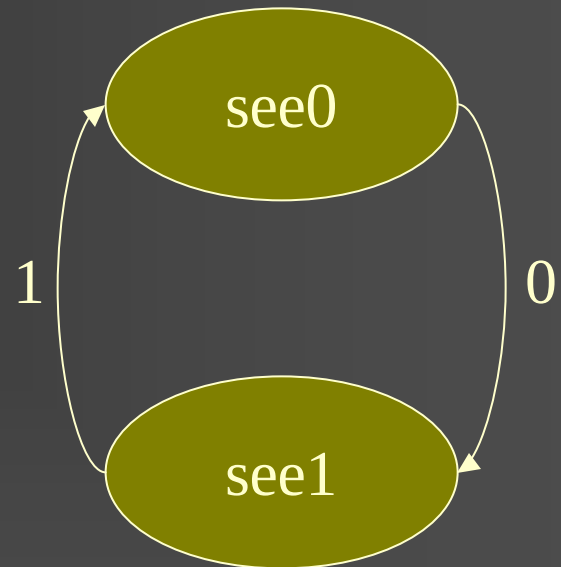
Let's Think About Automata

I want to be able to write

automaton see0

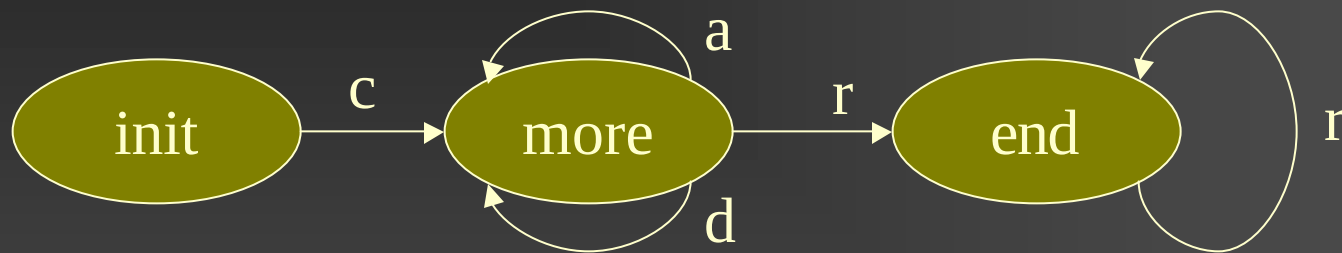
see0 : 0 $\rightarrow$ see1

see1 : 1 $\rightarrow$ see0



Another Example

car, cdr, cadr, cddr, cdar, caddr, ...



init : c $\rightarrow$ more

more : a $\rightarrow$ more

d $\rightarrow$ more

end : r $\rightarrow$ end

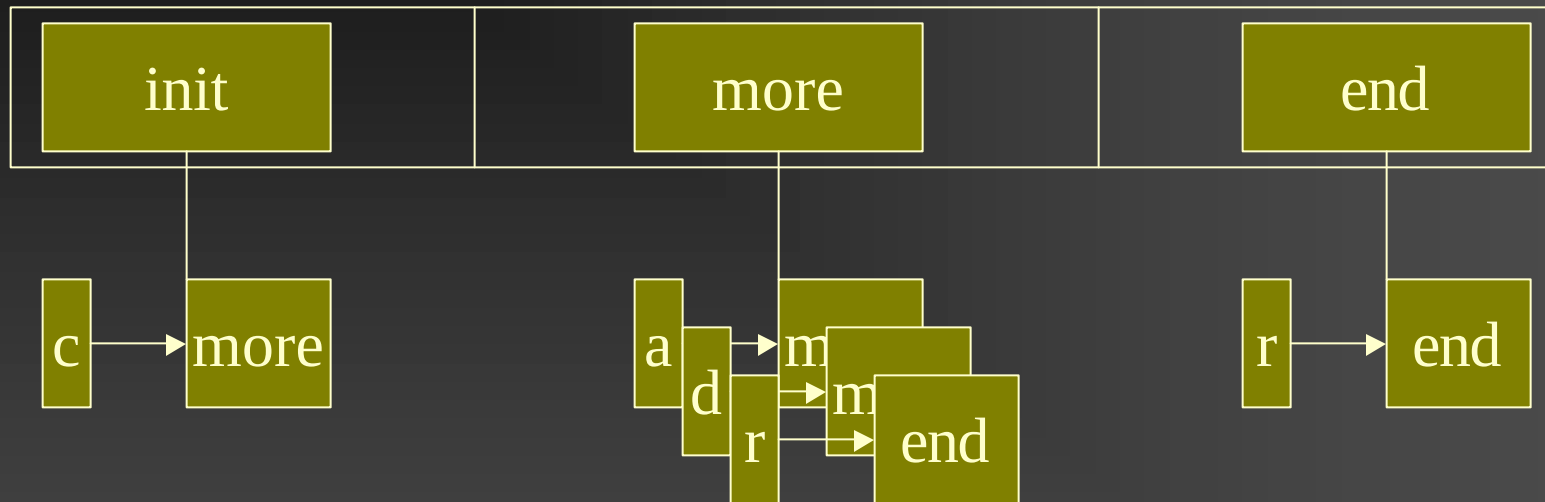
r $\rightarrow$ end

Let's Look at That Again

```
automaton init
  init : c → more
  more : a → more
         d → more
         r → end
  end   : r → end
```

How would **you** implement it?

Natural Solution



If stream ends, accept

If no next state found, reject

If next state found, continue

First Version

- Development time: 12 minutes
 - Bugs: 2
 - Performance (PLT Scheme):
 - 10000 elements: 50 ms
 - 100000 elements: 440 ms
 - 1000000 elements: 4316 ms
-

The Code

```
(define b-machine-states
  '((init (c more))
    (more (a more)
          (d more)
          (r end))
    (end (r end))))

(define (b-machine stream)
  (letrec ([walker (lambda (state stream)
                     (or (empty? stream)
                         (let ([transitions
                               (cdr (assv state b-machine-states))])
                           (let ([1st (first stream)])
                             (let ([new-state (assv 1st transitions)])
                               (if new-state
                                   (walker (cadr new-state) (rest stream))
                                   false)))))))]
    (walker 'init stream)))
```

What's The Essence?

- Per state, fast conditional dispatch table
 - An array of states
 - Quick state transition
-

A Message From Our Sponsors

We will encourage you to develop the three great virtues of a programmer: laziness, impatience, and hubris.

—Larry Wall and Randal L Schwartz

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Thinking Virtuously

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Compiler writers call this “case ... switch”
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Thinking Virtuously

- Per state, fast conditional dispatch table
Compiler writers call this “case ... switch”
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Function pointers offer random access
 - Quick state transition
If only function calls were implemented as “goto”s ...
-

In Other Words: `init` State Would Become

```
init ≡  
(procedure (stream)  
  (or (empty? stream)  
    (case (first stream)  
      [c (more (rest stream))]  
      [else false])))
```

In Other Words: `init` State Would Become

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    (case (first stream)  
      [c (more (rest stream))]  
      [else false]))))
```



more

In Other Words: more State Would Become

```
more ≡  
(procedure (stream)  
  (or (empty? stream)  
    (case (first stream)  
      [a (more (rest stream))]  
      [d (more (rest stream))]  
      [r (end (rest stream))]  
      [else false]))))
```

In Other Words: more State Would Become

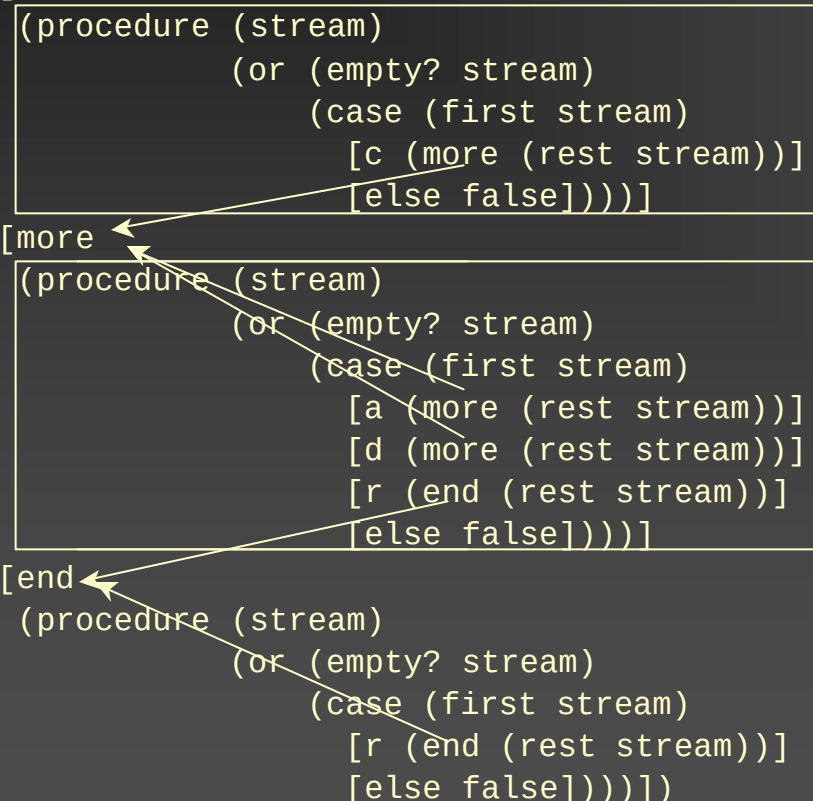
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more ≡  
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  (or (empty? stream)  
    (case (first stream)  
      [a (more (rest stream))]  
      [d (more (rest stream))]  
      [r (end (rest stream))]  
      [else false])))) → end
```

In Other Words: The Whole Code Would Become

```
(define b
  (letrec ([init
            (procedure (stream)
                      (or (empty? stream)
                          (case (first stream)
                            [c (more (rest stream))]
                            [else false]))))]
    [more
     (procedure (stream)
               (or (empty? stream)
                   (case (first stream)
                     [a (more (rest stream))]
                     [d (more (rest stream))]
                     [r (end (rest stream))]
                     [else false]))))]
    [end
     (procedure (stream)
               (or (empty? stream)
                   (case (first stream)
                     [r (end (rest stream))]
                     [else false]))))]
    init))
```

In Other Words: The Whole Code Would Become

```
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  (letrec ([init
            (procedure (stream)
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             [else false]))))]
    [end
     (procedure (stream)
       (or (empty? stream)
           (case (first stream)
             [r (end (rest stream))]
             [else false]))))]
    init))
```



The diagram illustrates the structure of the code. Three procedure objects are shown as boxes, each containing a definition. Arrows point from the labels [more], [end], and [init] to their respective procedure definitions. The [more] procedure is defined as a lambda function that takes a stream and returns either (more (rest stream)) if the first element is 'a' or 'd', or (end (rest stream)) if the first element is 'r', or false otherwise. The [end] procedure is defined as a lambda function that takes a stream and returns either (end (rest stream)) if the first element is 'r', or false otherwise. The [init] procedure is defined as a lambda function that takes a stream and returns either (more (rest stream)) if the first element is 'c', or false otherwise.

Scoreboard

- Laziness: 
 - Impatience: nope; too much code
 - Hubris: ...
-

In General

(state : (label → target) ...)

→

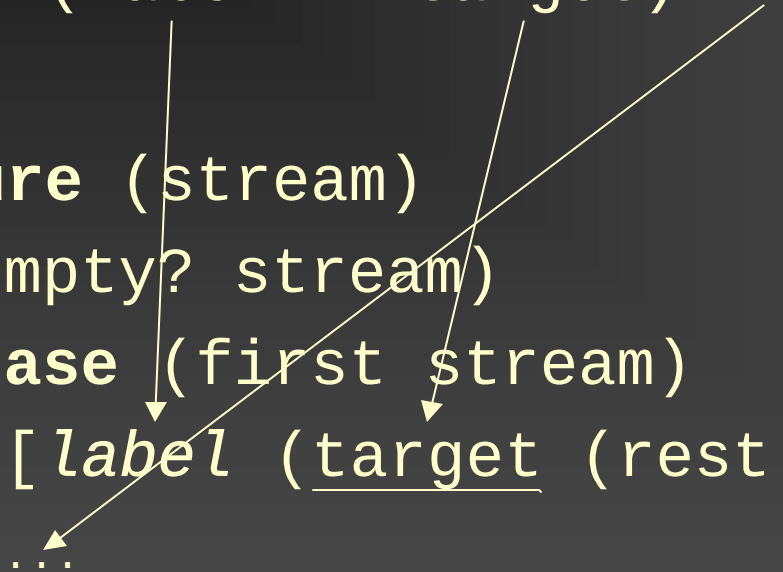
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In General

(state : (label → target) ...)

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```
(procedure (stream)
  (or (empty? stream)
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        [label (target (rest stream))]
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        [else false])))
```



Even More Generally

```
(_ init-state  
  (state : (cndn -> new-state) ...)  
  ...)
```



```
(letrec ([state  
  (procedure (stream)  
    (or (empty? stream)  
      (case (first stream)  
        [cndn (new-state (rest stream))]  
        ...  
        [else false]))))]  
  ...)  
  init-state)
```

Even More Generally

```
(_ init-state  
  (state : (cndn -> new-state) ...) )
```



```
(letrec ([state  
  (procedure (stream)  
    (or (empty? stream)  
        (case (first stream)  
          [cndn (new-state (rest stream))]  
          ...  
          [else false]))))]  
  ...)  
  init-state)
```

In Fact, That's the Code!

```
(define-syntax automaton
  (syntax-rules (-> :)
    [[input pattern
      output pattern] ] ) )
```

This is a Scheme macro

The Automaton

```
automaton init
```

```
  init : c → more
```

```
  more : a → more
```

```
        d → more
```

```
        r → end
```

```
  end  : r → end
```

In Scheme

```
(automaton init
  (init : (c → more))
  (more : (a → more)
          (d → more)
          (r → end))
  (end   : (r → end)))
```

What a Schemer really sees

```
(automaton init
  (init : (c → more))
  (more : (a → more)
         (d → more)
         (r → end))
  (end   : (r → end)))
```

With Clients

```
(define (v s)
  ((if (eq? (first s) 'c)
    (automaton init
      (init : (c -> loop))
      (loop : (a -> loop)
              (d -> loop)
              (r -> end))
      (end : (r -> end)))
    (automaton see0
      (see0 : (0 -> see1))
      (see1 : (1 -> see0)))
    s))
```


Second Version

- Development time: 5 minutes
 - Bugs: 0
 - Performance:
 - 10000 elements: 30 ms
 - 100000 elements: 310 ms
 - 1000000 elements: 3110 ms
-

Scoreboard

■ Laziness:



■ Impatience:



■ Hubris:

stay tuned

What Really Happened

The traditional implementation is an
Interpreter

The macro system implements a
Compiler

from Scheme++ to Scheme – and lets you
reuse the existing Scheme compiler

Macros

- Clean, convenient spec of automata
 - Permits *nested* ... – “pattern matching”
 - Easy to create domain-specific language
 - Each module can have different macros
-

Tail Calls

Ensures that

state transition =

goto =

loop for free!

Notice tail *recursion* isn't enough!

(Oh, and try generating loop code ...)

Stupid Parenthetical Syntax

```
(automaton see0  
  (see0 (0 -> see1))  
  (see1 (1 -> see0)))
```

is clearly ugly, evil, and an insidious plot
hatched by misbegotten academics

Smart Parenthetical Syntax

```
<automaton see0
  <state name="see0">
    <trn> <from> 0 </from>
      <to> see1 </to> </trn> </state>
  <state name="see1">
    <trn> <from> 1 </from>
      <to> see0 </to> </trn> </state>
</automaton>
```

is a hip, cool, great new idea

Python vs. Scheme (python.org)

- Standard object system
 - Regular expressions, Internet connectivity
 - Many builtin data types
 - One standard implementation
 - Relatively main-stream syntax
 - Main-stream control structures
-

Python vs. PLT Scheme (python.org)

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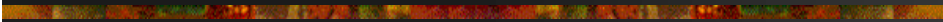
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 - ✓ ■ Many builtin data types
 - ✓ ■ One standard implementation
 - Relatively main-stream syntax – at what price?
 - Main-stream control structures
 - We got macros – you got five minutes?
 - Real Programmers use map/filter/fold, tail calls, ...
-

Take-Home Morals



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- If you claim to be smart, be *really* smart about reuse

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 - Scheme fits together particularly cleverly – you won't get it just by reading about it, you'll only *think* you did
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 - People who don't understand this use of tail calls don't get it
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Take-Home Morals

- If you claim to be smart, be *really* smart about reuse
 - Scheme fits together particularly cleverly – you won't get it just by reading about it, you'll only *think* you did
 - People who don't understand this use of tail calls don't get it
 - Take a real languages course in college
-

Scoreboard

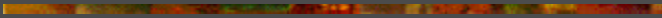
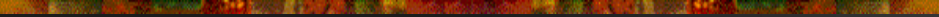
■ Laziness: 

■ Impatience: 

■ Hubris: 



A Parting Thought



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Obligatory URL

<http://www.plt-scheme.org/>

(Thanks: Matthias, Matthew, Robby, John, Paul, Paul,
Jamie, Philippe, Dorai, and dozens others)
