CSCI-1680 :: Computer Networks

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Based partly on lecture notes by David Mazières, Phil Levis, John Jannotti, Peterson & Davie

Cast

- Instructor: Rodrigo Fonseca (rfonseca)
- GTA: Andrew Ferguson (adf)
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Overview

- Goal: learn concepts underlying networks
 - How do networks work? What can one do with them?
 - Gain a basic understanding of the Internet
 - Gain experience writing *protocols*
 - Tools to understand new protocols and applications



Prerequisites

- CSCI-0320/CSCI-0360 (or equivalent).
 - We assume basic OS concepts (kernel/user, threads/ processes, I/O, scheduling)
- Low-level programming or be willing to learn quickly
 - threads, locking, explicit memory management, ...
- We allow any* language, but really *support* only C
 - You will be bit twiddling and byte packing...



Administrivia

- All assignments will be on the course page http://www.cs.brown.edu/courses/cs168/s11
- Text: Peterson and Davie, Computer Networks A Systems Approach, 4th Edition
- You are responsible to check the web page!
 - All announcements will be there
 - Textbook chapters corresponding to lectures: read them before class
 - Handouts, due dates, programming resources, *etc...*
 - *Subject to change* (reload before checking assignments)



Grading

- Exams: Midterm (15%) and Final (25%)
- Homework: Four written assignments (20%)
 - Short answer and design questions
- 4 Programming Projects (40%)
 - User level networking: streaming music server
 - IP, as an overlay, on top of UDP
 - TCP, on top of *your* IP
 - Final (TBD, we will solicit your input)



Networks

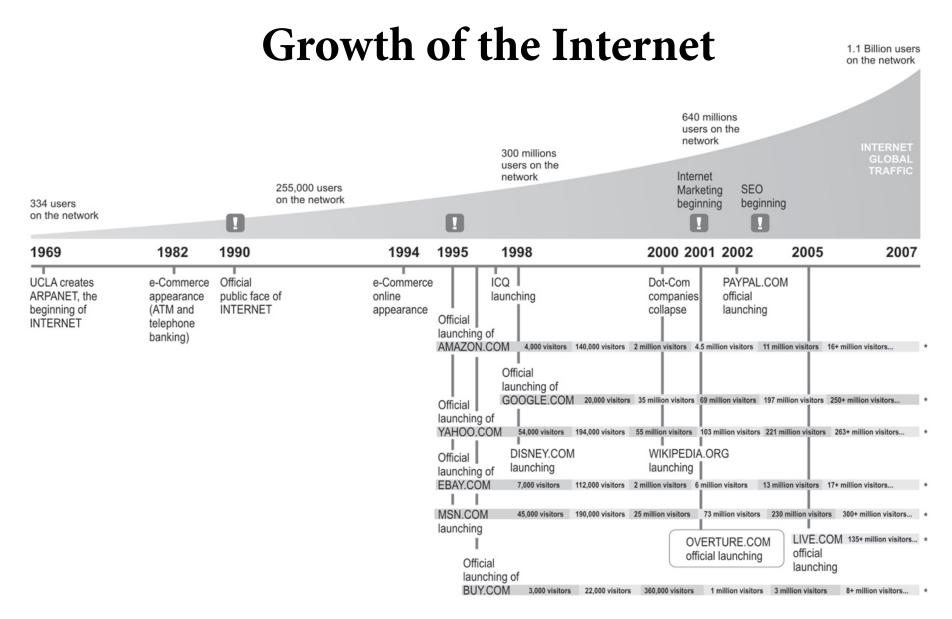
- What is a network?
 - System of lines/channels that interconnect
 - *E.g.*, railroad, highway, plumbing, postal, telephone, social, computer
- Computer Network
 - Moves information
 - Nodes: general-purpose computers (most nodes)
 - Links: wires, fiber optics, EM spectrum, composite...



Why Study Computer Networks?

- Many nodes are general-purpose computers
- Very easy to innovate and develop new uses of the network: *you* can program the nodes
- Contrast with the ossified Telephone network:
 - Can't program most phones
 - Intelligence in the network, control by parties vested in the *status quo*, ...





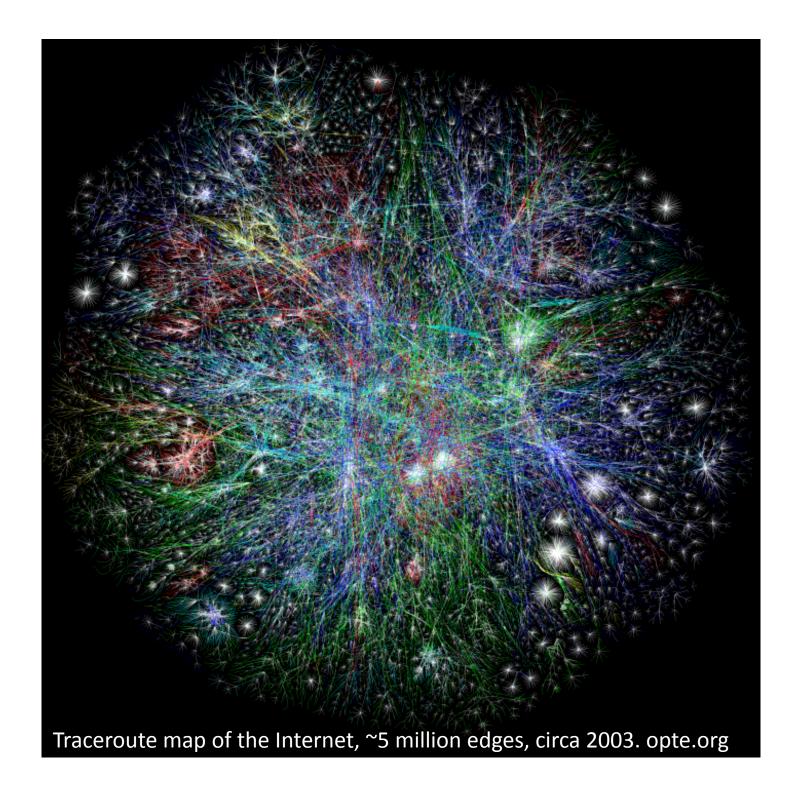


Source: Miguel Angel Todaro





Source: Facebook





Why should you take this course?

- Networks are cool!
 - Incredible impact: social, economic, political, educational, ...
- Incredible complexity
- Continuously changing and evolving
 - Any fact you learn will be inevitably out of date
 - Learn general underlying *principles*
- Learn to program the network



Roadmap

• Assignments: learn by implementing

- Warm up: Snowcast, a networked music server
 - Get a feel for how applications use the network

• Build knowledge from the ground up

- Link individual nodes
- Local networks with multiple nodes
- IP: Connect hosts across several networks
- Transport: Connect processes on different hosts
- Applications

• A few cross-cutting issues

- Security, multimedia, overlay networks, P2P...



Two-minutes for stretching

Building Blocks

- Nodes: Computers (hosts), dedicated routers, ...
- Links: Coax, twisted pair, fiber, radio, ...

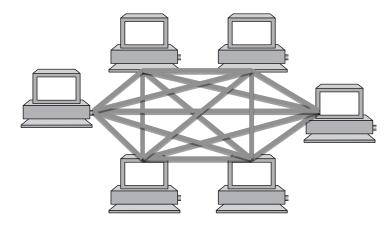


- Physical Layer: Several questions:
 - Voltage, frequency
 - Wired, wireless
- Link Layer: how to send data?
 - When to talk
 - What to say (format, "language")

Stay tuned for lectures 3 and 4...

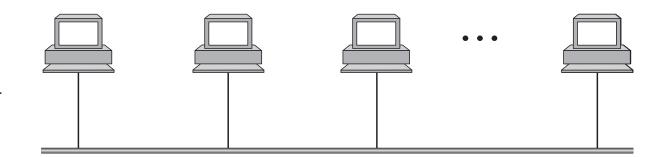


How to connect more nodes?



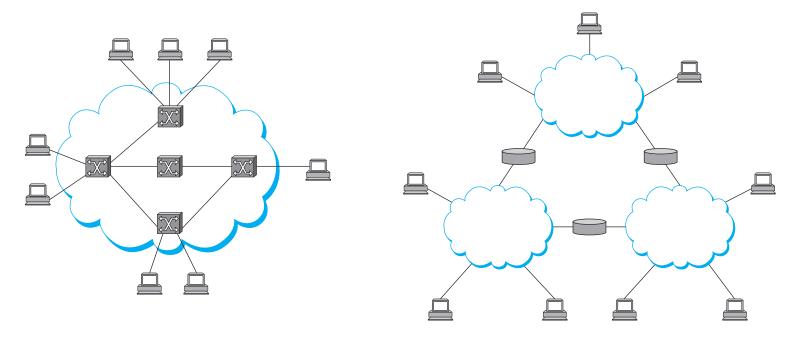
Multiple wires

Shared medium





From Links to Networks



- To scale to more nodes, use *switching*
 - Nodes can connect to multiple other nodes
 - Recursively, one node can connect to multiple networks



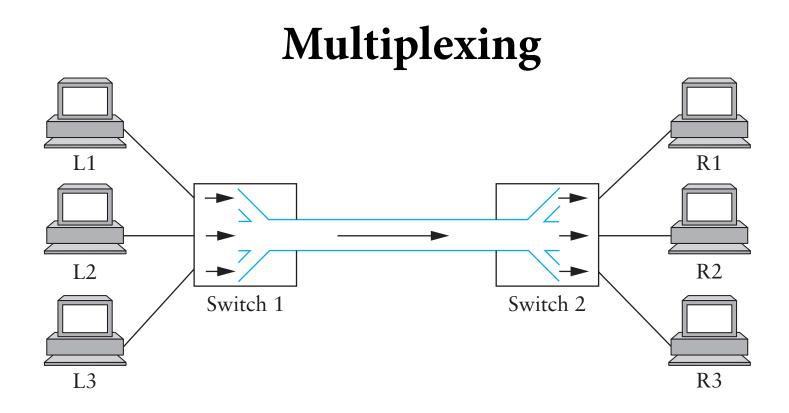
Switching Strategies

• Circuit Switching – virtual link between two nodes

- Set up circuit (*e.g.* dialing, signaling) may fail: busy
- Transfer data at known rate
- Tear down circuit
- Packet Switching
 - Forward bounded-size messages.
 - Each message can have different senders/receivers
 - Focus of this class

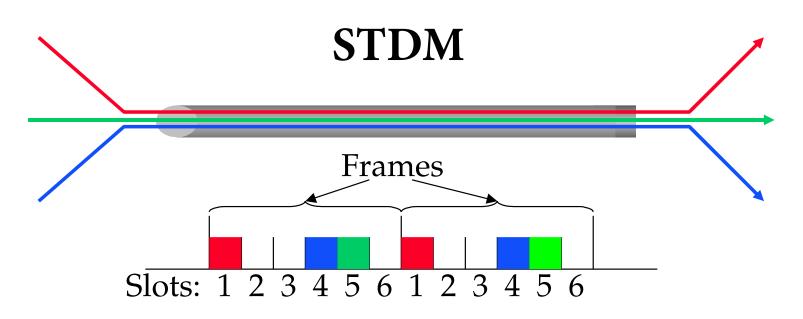
Analogy: circuit switching reserves the highway for a crosscountry trip. Packet switching interleaves everyone's cars.





• What to do when multiple flows must share a link?





Synchronous time-division multiplexing

- Divide time into equal-sized quanta, round robin
- Illusion of direct link for switched circuit net
- But wastes capacity if not enough flows
- Also doesn't degrade gracefully when more flows than slots

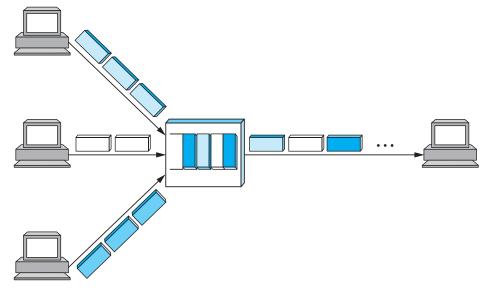


FDM

- Frequency-division multiplexing: allocates a frequency band for each flow
 - Same TV channels and radio stations
- Similar drawbacks to STDM
 - Wastes bandwidth if someone not sending
 - Can run out of spectrum



Statistical Multiplexing



- Idea: like STDM but with no pre-determined time slots (or order!)
- Maximizes link utilization
 - Link is never idle if there are packets to send



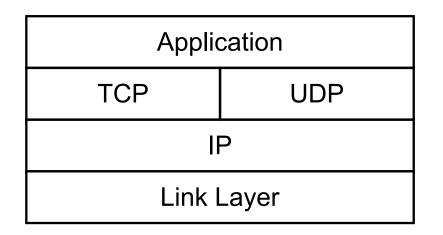
Statistical Multiplexing

• Cons:

- Hard to guarantee fairness
- Unpredictable queuing delays
- Packets may take different paths



Protocol Layering



- A network packet from A to D must be put in link packets A to B, B, to C, C to D
- Can view this encapsulation as a stack of layers
 - Each layer produces packets that become the payload of the lower layer's packets



End host End host Application Application Presentation Presentation Session Session Transport Transport Network Network Network Network Data link Data link Data link Data link Physical Physical Physical Physical One or more nodes within the network

OSI Reference Model



Layers

- Physical sends individual bits
- Data Link sends *frames*, handles media access
- Network sends packets, using routing
- Transport demultiplexes, provides reliability, flow and congestion control
- Session can tie together multiple streams (*e.g.*, audio & video)
- Presentation crypto, conversion between representations
- Application what the users sees, *e.g.*, HTTP



Addressing

- Each node typically has a unique* name
 - When that name also tells you how to get to the node, it is called an *address*
- Each layer can have its own naming/addressing
- *Routing* is the process of finding a path to the destination
 - In packet switched networks, each packet must have a destination address
 - For circuit switched, use address to set up circuit
- Special addresses can exist for broadcast/multicast/ anycast



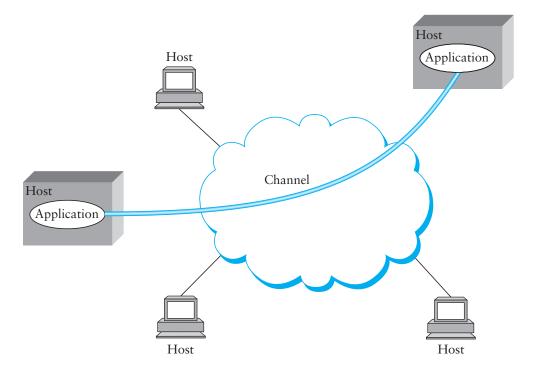
* or thinks it does, in case there is a shortage

Internet Protocol (IP)

- Used by most computer networks today
 - Runs *over* a variety of physical networks, can connect Ethernet, wireless, modem lines, etc.
- Every host has a unique 4-byte IP address (IPv4)
 - − *E.g.*, www.cs.brown.edu \rightarrow 128.148.32.110
 - The *network* knows how to route a packet to any address
- Need more to build something like the Web
 - Need naming (DNS)
 - Interface for browser and server software (next lecture)
 - Need demultiplexing within a host: which packets are for web browser, Skype, or the mail program?



Inter-process Communication



- Talking from host to host is great, but we want abstraction of inter-process communication
- Solution: encapsulate another protocol within IP



Transport: UDP and TCP

• UDP and TCP most popular protocols on IP

- Both use 16-bit *port* number & 32-bit IP address
- Applications *bind* a port & receive traffic on that port

• UDP – User (unreliable) Datagram Protocol

- Exposes packet-switched nature of Internet
- Sent packets may be dropped, reordered, even duplicated (but there is corruption protection)

• TCP – Transmission Control Protocol

- Provides illusion of reliable 'pipe' or 'stream' between two processes anywhere on the network
- Handles congestion and flow control

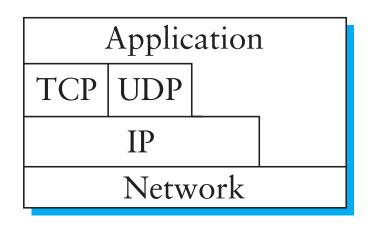


Uses of TCP

- Most applications use TCP
 - Easier to program (reliability is convenient)
 - Automatically avoids congestion (don't need to worry about taking down the network
- Servers typically listen on well-know ports:
 - SSH: 22
 - SMTP (email): 25
 - Finger: 79
 - HTTP (web): 80



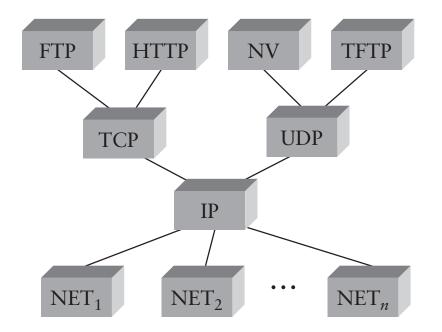
Internet Layering



- Strict layering not *required*
 - TCP/UDP "cheat" to detect certain errors in IP-level information like address
 - Overall, allows evolution, experimentation



IP as the Narrow Waist



- Many applications protocols on top of UDP & TCP
- IP works over many types of networks
- This is the "Hourglass" architecture of the Internet.
 - If every network supports IP, applications run over many different networks (*e.g.*, cellular network)



Coming Up

• Next class: how do applications use the network?

- Introduction to programming with Sockets
- Peterson & Davie 1.4
- Beej's Guide to Network Programming (link on the course website)
- Then...
 - We start our journey up the network stack, starting from how two computers can talk to each other.
- Remember: start your projects now!

