

Enhanced Approach Towards Network Reliability Using Boolean Algebra

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Motivation

- Computers are often connected to networks of various sizes
- Distributed Computer Systems (DCS) require reliable inter-computer communication
- Performance of such a network can be measured by determining the reliability of the overall network

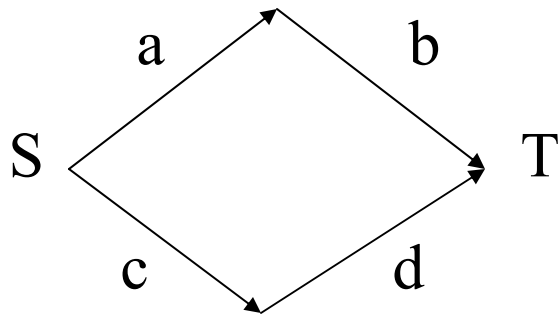
Problem

- The reliability analysis problem consists of determining the probability that a stochastic binary system operates
 - Example: a network for which each component/edge has an associated probability of working
- A coherent binary system is given by the set of minpaths P_i
 - Example: the set of minimal paths between two nodes whose functionality is sufficient for communication

Approach

- Use Boolean algebra, and in particular an SDP (sum of disjoint products) algorithm
- P_i can be viewed as a probabilistic event
 - Example: the probability that path P_i is working
- Make the events P_i disjoint and then add their corresponding probabilities to obtain the reliability of the entire system

Example of a simple network



- $P_1 = a b$, representing the event $a \wedge b$
- $P_2 = c d$, representing the event $c \wedge d$

Background

- A very simple method of making the events P_1, P_2, \dots, P_n disjoint:
 - $P_1, \neg P_1 P_2, \neg P_1 \neg P_2 P_3, \dots, \neg P_1 \neg P_2 \neg \dots \neg P_{n-1} P_n$,
where $\neg P_i$ means that P_i is failing
 - The order in which the minpaths P_i are listed can and does have an impact on the performance of any algorithm that is being used because this non-symmetrical pattern

Existing techniques

- SVI (single-variable inversions) [1]
 - to make two products disjoint, individual variables are negated
 - Eg: $\neg P_1 P_2 = \neg a c d + a \neg b c d$
- MVI (multiple-variable inversions) [2]
 - groups of variables can be negated together (disjunction of negations)
 - Eg: $\neg P_1 P_2 = \neg(a b) c d$

Proposed Approach

- GSDP (generalized sum of disjoint products)
 - Variables may be grouped together with any kind of Boolean expressions inside
 - No variable appears more than once in a product, to preserve the statistical independence of the groups
- Assumptions
 - Component states are mutually statistically independent
 - Each component of the network is either operational or failing

Example

$$P_1 = a b c x; \quad P_2 = a c d y; \quad P_3 = x y z$$

$$\text{SVI: } \neg P_1 \neg P_2 P_3 = \neg a x y z + a \neg c x y z + a c \neg b \neg d x y z$$

– 3 terms

$$\text{MVI: } \neg P_1 \neg P_2 P_3 = \neg(a c) x y z + a c \neg b \neg d x y z$$

– 2 terms

$$\text{GSDP: } \neg P_1 \neg P_2 P_3 = x y z (\neg(a c) \vee (\neg b \neg d))$$

– 1 term

Results and Discussion

- The Union of Products Problem (UPP), is NP-hard
- Two measures will be used to analyze the quality of the algorithm:
 - *computational efficiency*: practical run-time
 - resulting *number of disjoint products*: easy numerical evaluation and reduced computation and rounding errors.
- It is expected that this will reduce the size of the resulting reliability formula because several MVI terms can be combined into one GSDP term.

References

- [1] J.A. Abraham, “An Improved Algorithm for Network Reliability”, IEEE Trans. Reliability, vol R-28, no. 1 1979 Apr, pp 58-61.
- [2] K.D. Heidtmann, “Smaller Sums of Disjoint Products by Subproduct Inversion”, IEEE Trans. Reliability, vol. 38, no. 3, 1989 Aug, pp 305-311.