

CS 33

Data Representation, Part 1

Number Representation

- **Hindu-Arabic numerals**
 - developed by Hindus starting in 5th century
 - » positional notation
 - » symbol for 0
 - adopted and modified somewhat later by Arabs
 - » known by them as “Rakam Al-Hind” (Hindu numeral system)
 - **1999** rather than **MCMXCIX**
 - » (try doing long division with Roman numerals!)

Which Base?

- **1999**

- **base 10**

- » $9 \cdot 10^0 + 9 \cdot 10^1 + 9 \cdot 10^2 + 1 \cdot 10^3$

- **base 2**

- » **11111001111**

- $1 \cdot 2^0 + 1 \cdot 2^1 + 1 \cdot 2^2 + 1 \cdot 2^3 + 0 \cdot 2^4 + 0 \cdot 2^5 + 1 \cdot 2^6 + 1 \cdot 2^7 + 1 \cdot 2^8 + 1 \cdot 2^9 + 1 \cdot 2^{10}$

- **base 8**

- » **3717**

- $7 \cdot 8^0 + 1 \cdot 8^1 + 7 \cdot 8^2 + 3 \cdot 8^3$

- » **why are we interested?**

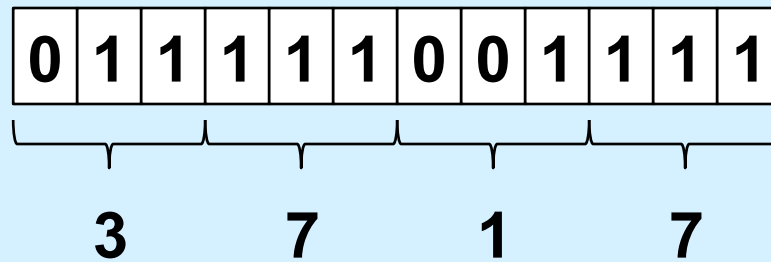
- **base 16**

- » **7CF**

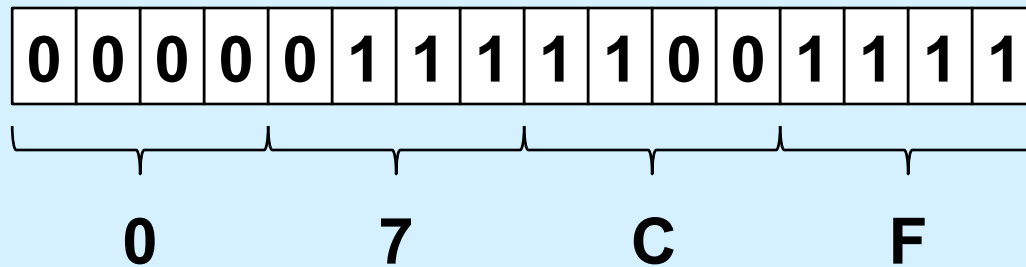
- $15 \cdot 16^0 + 12 \cdot 16^1 + 7 \cdot 16^2$

- » **why are we interested?**

Words ...



12-bit computer word



16-bit computer word

Algorithm ...

```
void baseX(unsigned int num, unsigned int base) {
    char digits[] = {'0', '1', '2', '3', '4', '5', '6', ... };
    char buf[8*sizeof(unsigned int)+1];
    int i;

    for (i = sizeof(buf) - 2; i >= 0; i--) {
        buf[i] = digits[num%base];
        num /= base;
        if (num == 0)
            break;
    }

    buf[sizeof(buf) - 1] = '\\0';
    printf("%s\\n", &buf[i]);
}
```

Or ...

```
$ bc
obase=16
1999
7CF
$
```

Quiz 1

- What's the decimal (base 10) equivalent of 23_{16} ?
 - a) 19
 - b) 33
 - c) 35
 - d) 37

Encoding Byte Values

- **Byte = 8 bits**
 - binary 00000000_2 to 11111111_2
 - decimal: 0_{10} to 255_{10}
 - hexadecimal 00_{16} to FF_{16}
 - » base 16 number representation
 - » use characters '0' to '9' and 'A' to 'F'
 - » write $FA1D37B_{16}$ in C as
 - $0xFA1D37B$
 - $0xfa1d37b$

Hex	Decimal	Binary
0	0	0000
1	1	0001
2	2	0010
3	3	0011
4	4	0100
5	5	0101
6	6	0110
7	7	0111
8	8	1000
9	9	1001
A	10	1010
B	11	1011
C	12	1100
D	13	1101
E	14	1110
F	15	1111

Boolean Algebra

- **Developed by George Boole in 19th Century**
 - algebraic representation of logic
 - » encode “true” as 1 and “false” as 0

And

- $A \& B = 1$ when both $A=1$ and $B=1$

$\&$	0	1
0	0	0
1	0	1

Or

- $A | B = 1$ when either $A=1$ or $B=1$

	0	1
0	0	1
1	1	1

Not

- $\sim A = 1$ when $A=0$

\sim	
0	1
1	0

Exclusive-Or (Xor)

- $A \wedge B = 1$ when either $A=1$ or $B=1$, but not both

\wedge	0	1
0	0	1
1	1	0

General Boolean Algebras

- Operate on bit vectors
 - operations applied bitwise

01101001	01101001	01101001	
<u>& 01010101</u>	<u> 01010101</u>	<u>^ 01010101</u>	<u>~ 01010101</u>
01000001	01111101	00111100	10101010

- All of the properties of boolean algebra apply

Example: Representing & Manipulating Sets

- Representation

- width- w bit vector represents subsets of $\{0, \dots, w-1\}$
- $a_j = 1$ iff $j \in A$

01101001 { 0, 3, 5, 6 }
76543210

01010101 { 0, 2, 4, 6 }
76543210

- Operations

&	intersection	01000001	{ 0, 6 }
	union	01111101	{ 0, 2, 3, 4, 5, 6 }
^	symmetric difference	00111100	{ 2, 3, 4, 5 }
~	complement	10101010	{ 1, 3, 5, 7 }

Bit-Level Operations in C

- **Operations &, |, ~, ^ available in C**
 - apply to any “integral” data type
 - » long, int, short, char
 - view arguments as bit vectors
 - arguments applied bit-wise
- **Examples (char datatype)**
 - $\sim 0x41 \rightarrow 0xBE$
 - $\sim 01000001_2 \rightarrow 10111110_2$
 - $\sim 0x00 \rightarrow 0xFF$
 - $\sim 00000000_2 \rightarrow 11111111_2$
 - $0x69 \& 0x55 \rightarrow 0x41$
 - $01101001_2 \& 01010101_2 \rightarrow 01000001_2$
 - $0x69 | 0x55 \rightarrow 0x7D$
 - $01101001_2 | 01010101_2 \rightarrow 01111101_2$

Contrast: Logic Operations in C

- **Contrast to Logical Operators**

- **&&, ||, !**

- » view 0 as “false”

- » anything nonzero as “true”

- » always return 0 or 1

- » early termination/short-circuited execution

- **Examples (char datatype)**

!0x41 → 0x00

!0x00 → 0x01

!!0x41 → 0x01

0x69 && 0x55 → 0x01

0x69 || 0x55 → 0x01

p && *p (avoids null pointer access)

Contrast: Logic Operations in C

- Contrast to Logical Operators

- &&, ||, !

- » view “false”

**Watch out for && vs. & (and || vs. |)...
One of the more common oopsies in
C programming**

- !0x41 → 0x00

- !0x00 → 0x01

- !!0x41 → 0x01

- 0x69 && 0x55 → 0x01

- 0x69 || 0x55 → 0x01

- p && *p (avoids null pointer access)

Shift Operations

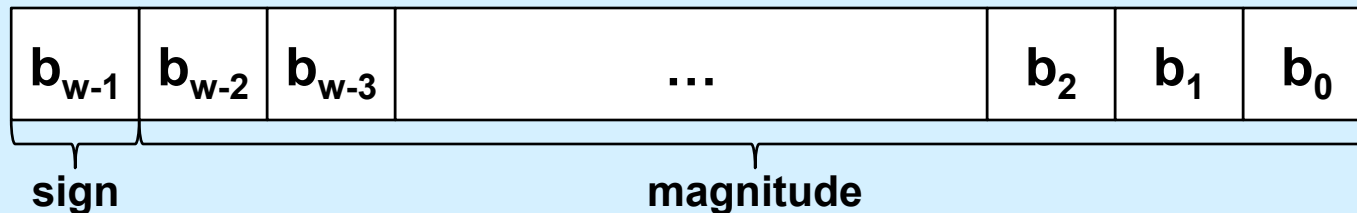
- **Left Shift:** $x \ll y$
 - shift bit-vector x left y positions
 - throw away extra bits on left
 - » fill with 0's on right
- **Right Shift:** $x \gg y$
 - shift bit-vector x right y positions
 - » throw away extra bits on right
 - **logical shift**
 - » fill with 0's on left
 - **arithmetic shift**
 - » replicate most significant bit on left
- **Undefined Behavior**
 - shift amount < 0 or \geq word size

Argument x	01100010
$\ll 3$	00010000
Log. $\gg 2$	00011000
Arith. $\gg 2$	00011000

Argument x	10100010
$\ll 3$	00010000
Log. $\gg 2$	00101000
Arith. $\gg 2$	11101000

Signed Integers

- **Sign-magnitude**



$$\text{value} = (-1)^{b_{w-1}} \cdot \sum_{i=0}^{w-2} b_i \cdot 2^i$$

- **two representations of zero!**
 - **computer must have two sets of instructions**
 - **one for signed arithmetic, one for unsigned**

Signed Integers

- **Ones' complement**
 - negate a number by forming its bitwise complement
 - » e.g., $(-1) \cdot 01101011 = 10010100$

$$\text{value} = -b_{w-1}(2^{w-1} - 1) + \sum_{i=0}^{w-2} b_i \cdot 2^i$$

$$= \sum_{i=0}^{w-2} b_i \cdot 2^i \quad \text{if } b_{w-1} = 0$$

$$= \sum_{i=0}^{w-2} (b_i - 1) \cdot 2^i \quad \text{if } b_{w-1} = 1$$

} two zeroes!

Signed Integers

- Two's complement

$b_{w-1} = 0 \Rightarrow$ non-negative number

$$\text{value} = \sum_{i=0}^{w-2} b_i \cdot 2^i$$

$b_{w-1} = 1 \Rightarrow$ negative number

$$\text{value} = (-1) \cdot 2^{w-1} + \sum_{i=0}^{w-2} b_i \cdot 2^i$$

one zero!

Signed Integers

- **Negating two's complement**

$$value = -b_{w-1} 2^{w-1} + \sum_{i=0}^{w-2} b_i 2^i$$

- **how to compute $-value$?**
 $(\sim value)+1$

Signed Integers

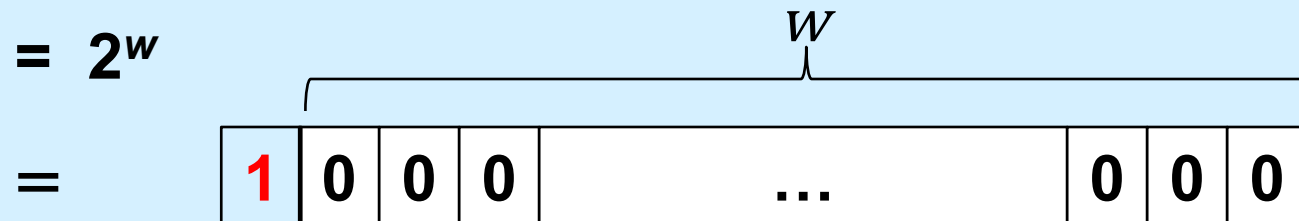
- Negating two's complement (continued)

$$value + (\sim value + 1)$$

$$= (value + \sim value) + 1$$

$$= (2^w - 1) + 1$$

$$= 2^w$$



Quiz 2

- **We have a computer with 4-bit words that uses two's complement to represent negative numbers. What is the result of subtracting 0010 (2) from 0001 (1)?**
 - a) 0111
 - b) 1001
 - c) 1110
 - d) 1111